# Special Graceful Labelings of Irregular Fences and Lobsters 

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#### Abstract

Irregular fences are subgraphs of $P_{m} \times P_{n}$ formed with $m$ copies of $P_{n}$ in such a way that two consecutive copies of $P_{n}$ are connected with one or two edges; if two edges are used, then they are located in levels separated an odd number of units. We prove here that any of these fences admits a special kind of graceful labeling, called $\alpha$-labeling. We show that there is a huge variety of this type of fences presenting a closed formula to determine the number of them that can be built on the grid $[1, m] \times[1, n]$. If only one edge is used to connect any pair of consecutive copies of $P_{n}$, the resulting graph is a tree. We use the $\alpha$-labelings of this type of fences to construct and label a subfamily of lobsters, partially answering the long standing conjecture of Bermond that states that all lobsters are graceful. The final labeling of the lobsters presented here is not only graceful, it is an $\alpha$-labeling, therefore they can be used to produce new graceful trees.


## 1. Introduction

Suppose $G$ is a graph of order $n$ and size $m$. An injective function $f: V(G) \rightarrow\{0,1, \ldots, m\}$ is called a graceful labeling of $G$ if every edge $u v$ of $G$ has assigned a weight, defined by $|f(u)-f(v)|$, and the set of all weights induced by $f$ on the edges of $G$ is $\{1,2, \ldots, m\}$. A graph that admits a graceful labeling is called graceful. This labeling, together with three other labelings, was introduced by Rosa [1] as a mean to study a problem in combinatorial design associated with the decomposition of the complete graph $K_{2 m+1}$ into copies of any tree of size $m$. Rosa proved that if there is a graceful labeling of a tree $T$ of size $m$, then there exists a (cyclic) decomposition of $K_{2 m+1}$ into copies of $T$. Several applications of gracefully labeled graphs are known, we can mentioned here the ones presented by Bloom and Golomb [2] and [3], and the ones given by Brankovic and Wanless [4].

An $\alpha$-labeling of $G$ is a graceful labeling $f$ for which there exists an integer $\lambda$, called the boundary value of $f$, such that for each edge $u v$ of $G$, either $f(u) \leq \lambda<f(v)$ or $f(v) \leq \lambda<f(u)$. If $G$ admists an $\alpha$-labeling, then it is called an $\alpha$-graph. This definition of an $\alpha$-graph implies that $G$ is bipartite and $\lambda$ is the smaller of the two vertex labels that yield the weight 1 . This type of labeling is the most restrictive one among the four labelings introduced by Rosa [1]. The existence of an $\alpha$-labeling implies the existence of several other types of labelings; so, they are located at the center of this research area. Not all graphs are graceful or $\alpha$, this fact motivates the search of new families of graphs admitting these types of labelings.

Let $G$ be a graph of order $n$ and size $m$. Suppose that $f$ is a graceful labeling of $G$. The labeling $\bar{f}: V(G) \rightarrow\{0,1, \ldots, m\}$, defined as $\bar{f}=m-f(v)$ for every $v \in V(G)$, is called the complementary labeling of $f$; this is also a graceful labeling; thus, its existence can be used to prove that the number of graceful labelings of any graph is always even. Let $g$ be a labeling of $G$ defined as $g(v)=c+f(v)$ for every $v \in V(G)$; we say that $g$ is a $c$ units shifting of $f$. It is not difficult to see that both, $f$ and $g$, induce the same weights. Suppose now that $f$ is an $\alpha$-labeling of $G$ with boundary value $\lambda$; the labeling $h$, defined for every $v \in V(G)$ as

$$
h(v)= \begin{cases}f(v) & \text { if } f(v) \leq \lambda \\ d-1+f(v) & \text { if } f(v)>\lambda\end{cases}
$$

is called a d-graceful labeling of $G$. This type of labeling was introduced in 1982, independently, by Maheo and Thuillier [5] and Slater [6].
Suppose that $f(v)-f(u)=w>0$, then $h(v)-h(u)=d-1+f(v)-f(u)=d-1+w$. Since $1 \leq w \leq m$, we get that $d \leq d-1+w \leq d-1+m$. In other terms, the weights induced by $h$ on the edges of $G$ are $d, d+1, \ldots, d-1+m$. This property of the $\alpha$-labelings has been widely used to construct new graceful and $\alpha$-graphs starting with smaller $\alpha$-graphs. The reverse of $f$, denoted by $f_{r}$, is another $\alpha$-labeling of $G$, it is defined as

$$
f_{r}(v)= \begin{cases}\lambda-f(v) & \text { if } f(v) \leq \lambda \\ m+\lambda+1-f(v) & \text { if } f(v)>\lambda\end{cases}
$$

Note that $f$ and $f_{r}$ have the same boundary value; in addition, if $f(v)-f(u)=w$, for any weight $w \in\{1,2, \ldots, m\}$, then $f_{r}(v)-f_{r}(u)=$ $m+\lambda+1-f(v)-\lambda+f(u)=m+1-(f(v)-f(u))=m+1-w$.

In Section 2 we present an $\alpha$-labeling for a large family of connected subgraphs of the grid $P_{m} \times P_{n}$. This family, denoted by $\mathscr{F}$, is formed by all the graphs built in the following way:

For every $i \in\{1,2, \ldots, m\}$, let $P_{n}^{i}$ be the path of order $n$ with vertex set $V\left(P_{n}^{i}\right)=\left\{v_{i, 0}, v_{i, 1}, \ldots, v_{i, n-1}\right\}$ and edge set $E\left(P_{n}^{i}\right)=\left\{v_{i, 0} v_{i, 1}\right.$, $\left.v_{i, 1} v_{i, 2}, \ldots, v_{i, n-2} v_{i, n-1}\right\}$. Now, for every $i \in\{1,2, \ldots, m-1\}$, decide whether $P_{n}^{i}$ is connected to $P_{n}^{i+1}$ with one or two edges (also called links). If only one edge connects them, then choose any $j \in\{0,1, \ldots, n-1\}$ and connect with an edge the vertices $v_{i, j}$ and $v_{i+1, j}$. If two edges connect them, then choose $j_{1}, j_{2} \in\{0,1, \ldots, n-1\}$, where $\left|j_{2}-j_{1}\right|$ is odd, and introduce the edges $v_{i, j_{1}} v_{i+1, j_{1}}$ and $v_{i, j_{2}} v_{i+1, j_{2}}$. Given that the number of edges connecting two copies of $P_{n}$ may vary, we refer to this type of graph as an irregular fence. In Figure 1.1 we show all the nonisomorphic fences in $\mathscr{F}$ built on $[1,3] \times[1,4]$. We claim that all the irregular fences are $\alpha$-graphs.


Figure 1.1: All nonisomorphic irregular fences built on $[1,3] \times[1,4]$

In Section 3 we study this type of irregular fences from an enumerative perspective. We present a closed formula for the number of nonisomorphic irregular fences built on $[1, m] \times[1, n]$. When every pair of consecutive copies of $P_{n}$ is connected with only one edge, the resulting fence corresponds to a type of tree called path-like tree; it is known that they are $\alpha$-trees [7]. In Section 4 we consider a subfamily of the path-like trees built on $[1, m] \times[1,5]$, with the extra property that they are lobsters. We characterize the lobsters that are irregular fences, therefore, $\alpha$-trees; in addition we show that some other $\alpha$-lobsters can be obtained from them by adding pendant vertices to some or all the vertices at distance one from the central path.

All graphs considered in this work are simple, i.e., no loops nor multiple edges are allowed. We mainly follow the notation and terminology used in [8] and [9].

## 2. $\alpha$-labelings of irregular fences

As we mentioned before, $\alpha$-labelings were introduced by Rosa [1]; he presented a labeling scheme for caterpillars that can be easily adapted for the case of paths. For the sake of completeness, we present here Rosa's $\alpha$-labeling of the path $P_{n}$; we use this labeling in the construction of the $\alpha$-labeled irregular fences.
Lemma 2.1. For every $n \geq 1$, the path $P_{n}$ is an $\alpha$-graph.
Assuming that $V\left(P_{n}\right)=\left\{v_{0}, v_{1}, \ldots, v_{n-1}\right\}$ and $E\left(P_{n}\right)=\left\{v_{0} v_{1}, v_{1} v_{2}, \ldots, v_{n-2} v_{n-1}\right\}$, the $\alpha$-labeling $f: V\left(P_{n}\right) \rightarrow\{0,1, \ldots, n-1\}$ is defined as:

$$
f\left(v_{i}\right)= \begin{cases}\frac{i}{2} & \text { if } i \text { is even } \\ n-\frac{i+1}{2} & \text { if } i \text { is odd }\end{cases}
$$

The labeling $f$ has boundary value $\lambda=\frac{n-2}{2}$ when $n$ is even and $\lambda=\frac{n-1}{2}$ when $n$ is odd. Moreover, $f\left(v_{0}\right)=0$ regardless the parity of $n$ but $f\left(v_{n-1}\right)=\frac{n}{2}=\lambda+1$ when $n$ is even and $f\left(v_{n-1}\right)=\frac{n-1}{2}=\lambda$ when $n$ is odd. We say that $v \in V\left(P_{n}\right)$ is a black vertex if $f(v) \leq \lambda$, otherwise $v$ is a white vertex. In Figure 2.1 we show two examples of this labeling on $P_{12}$ and $P_{17}$. Just for the examples, the boundary value is on a red vertex.


Figure 2.1: $\alpha$-labelings of $P_{12}$ and $P_{17}$

The construction of the $\alpha$-labeled irregular fences, built on $P_{m} \times P_{n}$, is based on an embedding of the path $P_{m n}$ on the grid $[1, m] \times[1, n]$. The division algorithm tell us that for each $i \in\{1,2, \ldots, m n\}$, there exist unique $q$ and $r$ such that $i=q n+r$, where $0 \leq r<n$. Using this fact we can define the embedding of $P_{m n}$ on the grid $[1, m] \times[1, n]$ to be the bijective function $\phi:\left\{v_{0}, v_{1}, \ldots, v_{m n-1}\right\} \rightarrow[1, m] \times[1, n]$, where

$$
\phi\left(v_{i}\right)= \begin{cases}(q+1, r+1) & \text { if } q \text { is even }, \\ (q+1, n-r) & \text { if } q \text { is odd. }\end{cases}
$$

Once the embedding is done, we proceed to label the vertices of $P_{m n}$ using the function $f$ given in Lemma 2.1. In the first part of Figure 2.2 we show an embedding of $P_{15}$ on the grid $[1,5] \times[1,3]$, on the second part we exhibit the $\alpha$-labeling of this path at this embedding.


Figure 2.2: Embedding of $P_{15}$ on $[1,5] \times[1,3]$ and its $\alpha$-labeling

In the following lemmas we present the essential results that will allow us to prove that any irregular fence in $\mathscr{F}$ is an $\alpha$-graph.
Lemma 2.2. Any fence built on $[1,2] \times[1, n]$, with only one edge connecting the two copies of $P_{n}$, is an $\alpha$-graph.
Proof. Suppose that $P_{2 n}$ has been embedded in the grid $[1,2] \times[1, n]$ in the way described before. In addition, assume that $P_{2 n}$ has been labeled using the function $f$ given in Lemma 2.1. In the following diagram we show this labeling where the labels on the black vertices are at most $\lambda$, the boundary value of $f$, while the labels on the white vertices are at least $\lambda+1$. Note that the edge connecting the vertices on $(1, n)$ and $(2, n)$ has weight $y-x-5$, independently of the parity of $n$.


If for any feasible value of $t$, the vertices on $(1, n-t)$ and $(2, n-t)$ are connected, the new edge also has weight $y-x-5$. This implies that all the horizontal "edges" on this embedding of $P_{2 n}$ have the same weight and any of them can be used to connect the two copies of $P_{n}$, being the final fence an $\alpha$-graph.

Lemma 2.3. Any fence built on $[1,2] \times[1, n]$, with two edges connecting the two copies of $P_{n}$, is an $\alpha$-graph.
Proof. As we did in Lemma 2.2, suppose that $P_{2 n}$ has been embedded in the grid $[1,2] \times[1, n]$, in the way described before, and that it has been labeled using the $\alpha$-labeling $f$ in Lemma 2.1. In the following diagram, we show new labelings for the two copies of $P_{n}$.


These labelings are obtained from $f$ by fixing the labels on the black vertices of the first copy of $P_{n}$ and adding one unit to all other vertices. In this way, the edges on the first copy of $P_{n}$ have the weights $n+2, n+3, \ldots, 2 n$; the weights on the edges of the second copy of $P_{n}$ are $1,2, \ldots, n-1$. We use all the labels in $\{0,1, \ldots, 2 n\}$ except $\left\lceil\frac{n}{2}\right\rceil$. Since the white vertices on the second copy of $P_{n}$ were augmented one unit while the black vertices on the first copy were fixed, any line connecting a black vertex with a white vertex will be an edge of weight $y-x-4$. Similarly, any line connecting a white vertex with a black vertex will be an edge of weight $y-x-5$ because the labels of both endvertices were augmented one unit. Hence, by connecting both copies of $P_{n}$ with two edges, one of each kind, that is, one black-white and one white-black, we obtain an $\alpha$-labeled irregular fence. This fence is in $\mathscr{F}$ because these types of edges are in alternated levels. This concludes the proof.

In Figure 2.3 we show four examples of these labeled irregular fences, two for each lemma.


Figure 2.3: $\alpha$-labelings of four irregular fences

Theorem 2.4. If $G$ is an irregular fence in $\mathscr{F}$, then $G$ is an $\alpha$-graph.
Proof. Suppose that $G$ is an irregular fence built on $P_{m} \times P_{n}$ such that it contains $1 \leq k \leq m-1$ pairs of consecutive copies of $P_{n}$ connected by two edges. Thus, $G$ has size $m(n-1)+(m-1)+k=m n-1+k$. Assume that the path $P_{m n}$ has been labeled using the function $f$ in Lemma 2.1 and is embedded in the grid $[1, m] \times[1, n]$. Thus, the weights induced on the edges of every copy of $P_{n}$ are consecutive integers, and the horizontal edges, of this embedding of $P_{m n}$, have weights $(m-1) n,(m-2) n, \ldots, 2 n, n$.
Now we delete all the horizontal edges connecting consecutive copies of $P_{n}$ in $P_{m n}$. Once this is done, we draw new horizontal edges following the pattern in $G$, In this way, we have a labeling of $G$; based on Lemma 2.2, this is an $\alpha$-labeling when $G$ is a tree, that is, when only one edge connects any pair of consecutive copies of $P_{n}$. If this is not the case, i.e., when there are $k>0$ pairs of consecutive copies of $P_{n}$ connected with two edges, these two horizontal edges have the same weight. To eliminate this duplicity, we apply the procedure used in the proof of Lemma 2.3.
Suppose that $i_{1}, i_{2}, \ldots, i_{k}$ are the indices for which there are two horizontal edges connecting $P_{n}^{i_{j}}$ and $P_{n}^{i_{j}+1}$. For every $i \leq i_{j}$, the labels of the black vertices of all $P_{n}^{i}$ are fixed and all the other labels are augmented in one unit. In this way, these horizontal edges have different
weights that are consecutive integers. Once this process has been applied to every pair of consecutive copies of $P_{n}$ connected by two edges, the resulting labeling is indeed an $\alpha$-labeling of $G$. In fact, since there are exactly $k$ pairs of consecutive copies of $P_{n}$ connected by two edges, the original labels of the white vertices have been shifted $k$ units, avoiding the duplicity of vertex labels; the weights on each copy of $P_{n}$ are consecutive integers, and the weights on the horizontal edges complement the ones on the vertical edges. Therefore, the final labeling of $G$ is an $\alpha$-labeling and $G$ is an $\alpha$-graph.

In Figure 2.4 we show an example of this labeling where $G$ is built on $P_{10} \times P_{10}$ and $k=7$.


Figure 2.4: $\alpha$-labelings of a fence of size 106 built on $[1,10] \times[1,10]$

## 3. Enumerating irregular fences

Motivated by the result in the previous section, we want to determine the number of this type of fences. In [10], we found the number of fences that can be built on the grid $[1, m] \times[1, n]$. Using that result, we present here a closed formula for the number of nonisomorphic irregular fences built on the grid.
We start by counting the number of irregular fences that can be built on $[1,2] \times[1, n]$. Since the grid $[1, m] \times[1, n]$ can be seen as a linear amalgamation of $m-1$ copies of $[1,2] \times[1, n]$ we refer to the fences on $[1,2] \times[1, n]$ as building blocks, or just blocks, of $[1, m] \times[1, n]$. Thus, a block in an irregular fence consists of two copies of $P_{n}$ and 1 or 2 (horizontal) links (edges), It is not difficult to see that the number of blocks with only one link is $C(n, 1)=n$, i.e., the number of ways of selecting one element from $\{1,2, \ldots, n\}$. To determine the number of blocks with two links we may count the 2-element subsets of $\{1,2, \ldots, n\}$, such that the difference between the two elements is odd. Thus, for any subset $\{i, j\}$, with $i<j$, the possible values for $j$ are determined by the value of $i$. When $i$ is odd, there are $\left\lfloor\frac{n}{2}\right\rfloor-\frac{i-1}{2}$ possible values for $j$. When $i$ is even, there are $\left\lceil\frac{n}{2}\right\rceil-\frac{i}{2}$ possible values for $j$.
Hence, when $n$ is even, the number of 2-element subsets satisfying the conditions is given by

$$
\sum_{i=1}^{\frac{n}{2}} i+\sum_{i=1}^{\frac{n}{2}-1} i=2 \sum_{i=1}^{\frac{n}{2}-1} i+\frac{n}{2}=\frac{2\left(\frac{n}{2}-1\right) \frac{n}{2}}{2}+\frac{n}{2}=\frac{n}{2}\left(\frac{n}{2}-1+1\right)=\frac{n^{2}}{4}
$$

When $n$ is odd, this number is

$$
2 \sum_{i=1}^{\frac{n-1}{2}} i=\frac{2\left(\frac{n-1}{2}\right)\left(\frac{n+1}{2}\right)}{2}=\frac{n^{2}-1}{4}
$$

Therefore, the number of blocks is $n+\frac{n^{2}}{4}=\frac{n^{2}+4 n}{4}$ when $n$ is even and $n+\frac{n^{2}-1}{4}=\frac{n^{2}+4 n-1}{4}$ when $n$ is odd. For $n \geq 1$, the sequence $a(n)$ formed by these values corresponds to the sequence A002620 in OEIS [11].

Another number needed in our counting process is the number of symmetric blocks. Once again, we start analyzing the case where the block has exactly one link. If $n$ is even, there are no symmetric blocks. If $n$ is odd, there is only one symmetric block. We have a similar situation when the block has two links. When $n$ is odd there are no two numbers $i<j$ in $\{1,2, \ldots, n\}$ such that $j-i$ is odd and $i-1=n-j$. When $n$ is even, for every $1 \leq i \leq \frac{n}{2}$, the number $j=n+1-i$ belongs to $\left\{\frac{n}{2}+1, \frac{n}{2}+2, \ldots, n\right\}, j-i=n+1-i-i=n+1-2 i$ is odd and $i-1=n-j=n-(n+1-i)=i-1$. Then, if $s(n)$ denotes the number of symmetric blocks, we get

$$
s(n)= \begin{cases}\frac{n}{2} & \text { if } n \text { is even } \\ 1 & \text { if } n \text { is odd }\end{cases}
$$

For $n \geq 1$, the sequence $s(n)$ is the sequence A152271 in OEIS [12].
Now we turn our attention to the general case. Given any irregular fence $F$ built on $[1, m] \times[1, n]$, there are other three fences that are isomorphic to $F$ : when $F$ is rotated $180^{\circ}$ around a central vertical axis, when $F$ is rotated $180^{\circ}$ around a central horizontal axis, and when $F$ is rotated $180^{\circ}$ around a central axis perpendicular to the plane containing $F$. Thus, there are three possible situations: $F$ has four different representations, $F$ has two different representations, or $F$ has one representation. Let $T$ be the set of all irregular fences on $[1, m] \times[1, n]$; we define $V$ to be the subset of $T$ containing the fences with a vertical symmetry, $H$ to be the subset of $T$ containing the fences with a horizontal symmetry, $C$ to be the subset of $T$ containing the fences with a central symmetry, and $A$ to be the subset of $T$ containing the fences with all these symmetries. In Figure 3.1 we show four examples, one for each of these subsets.


Figure 3.1: Different types of symmetric fences

Since the fences in $A$ have all the described symmetries, each of them appears only once in the list of all possible fences built on $[1, m] \times[1, n]$. Every element of $V-A, H-A$, or $C-A$ appears twice in this list. Every nonsymmetric fence appears four times in the list. Thus, if we take the addition of cardinalities

$$
|T|+|V|+|H|+|C|
$$

every fence is counted four times. Therefore, the number of nonisomorphic irregular fences built on $[1, m] \times[1, n]$ is given by

$$
f(m, n)=\frac{1}{4}(|T|+|V|+|H|+|C|)
$$

In order to find a closed formula for $f(m, n)$ we just need to determine explicitely these four cardinalities.
Based on the number of blocks and symmetric blocks, found above, and the fact that $[1, m] \times[1, n]$ can be formed with $m-1$ copies of $[1,2] \times[1, n]$, we can say that

$$
|T|= \begin{cases}\left(\frac{n^{2}+4 n}{4}\right)^{m-1} & \text { if } n \text { is even } \\ \left(\frac{n^{2}+4 n-1}{4}\right)^{m-1} & \text { if } n \text { is odd }\end{cases}
$$

If $F$ is a fence in $V$, then its $i$ th block is identical to its $(m-i)$ th block. This implies that we need to determine the number of posibilities for the first $\left\lfloor\frac{m}{2}\right\rfloor$ blocks. Thus,

$$
|V|= \begin{cases}\left(\frac{n^{2}+4 n}{4}\right)^{\left\lfloor\frac{m}{2}\right\rfloor} & \text { if } n \text { is even } \\ \left(\frac{n^{2}+4 n-1}{4}\right)^{\left\lfloor\frac{m}{2}\right\rfloor} & \text { if } n \text { is odd }\end{cases}
$$

If $F \in H$, each block in $F$ must be symmetric. So,

$$
|H|= \begin{cases}\left(\frac{n}{2}\right)^{m-1} & \text { if } n \text { is even } \\ 1 & \text { if } n \text { is odd }\end{cases}
$$

When $F \in C$, there are two cases that we need to analyze that depend on the parity of $m$. Recall that in this case the $i$ th block of $F$ is represented up side down in the $(m-i)$ th block.

If $m$ is even and $i=\frac{m}{2}$, then $i=m-i$. This implies that the $i$ th block of $F$ must be symmetric. So,

$$
|C|= \begin{cases}\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m-2}{2}} \cdot \frac{n}{2} & \text { if } n \text { is even } \\ \left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m-2}{2}} \cdot 1 & \text { if } n \text { is odd }\end{cases}
$$

If $m$ is odd

$$
|C|= \begin{cases}\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m-1}{2}} & \text { if } n \text { is even } \\ \left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m-1}{2}} & \text { if } n \text { is odd }\end{cases}
$$

Thus, we have found a closed formula for $F(m, n)$. We summarize these results in the following theorem.

Theorem 3.1. The number $f(m, n)$ of nonisomorphic irregular fences built on $[1, m] \times[1, n]$ is:

- When both $m$ and $n$ are even.
$f(m, n)=\frac{1}{4}\left(\left(\frac{n^{2}+4 n}{4}\right)^{m-1}+\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m}{2}}+\left(\frac{n}{2}\right)^{m-1}+\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m-2}{2}} \cdot \frac{n}{2}\right)$
- When $m$ is even and $n$ is odd.
$f(m, n)=\frac{1}{4}\left(\left(\frac{n^{2}+4 n-1}{4}\right)^{m-1}+\left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m}{2}}+1+\left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m-2}{2}}\right)$
- When $m$ is odd and $n$ is even. $f(m, n)=\frac{1}{4}\left(\left(\frac{n^{2}+4 n}{4}\right)^{m-1}+\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m-1}{2}}+\left(\frac{n}{2}\right)^{m-1}+\left(\frac{n^{2}+4 n}{4}\right)^{\frac{m-1}{2}}\right)$
- When both $m$ and $n$ are odd
$f(m, n)=\frac{1}{4}\left(\left(\frac{n^{2}+4 n-1}{4}\right)^{m-1}+\left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m-1}{2}}+1+\left(\frac{n^{2}+4 n-1}{4}\right)^{\frac{m-1}{2}}\right)$
In Table 1, read by rows, we show the first values of $f(m, n)$ for $2 \leq m, n \leq 10$. We have omitted the cases where $m=1$ or $n=1$ because $f(m, n)=1$.

|  | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 2 | 2 | 3 | 5 | 6 | 9 | 10 | 14 | 15 | 20 |
| 3 | 4 | 9 | 21 | 36 | 66 | 100 | 160 | 225 | 330 |
| 4 | 10 | 39 | 150 | 366 | 918 | 1810 | 3640 | 6315 | 11100 |
| 5 | 25 | 169 | 1060 | 3721 | 12789 | 32761 | 83296 | 177241 | 375925 |
| 6 | 70 | 819 | 8360 | 40626 | 190917 | 620830 | 1994944 | 5134095 | 13143500 |
| 7 | 196 | 3969 | 65808 | 443556 | 2849526 | 11764900 | 47783680 | 148718025 | 459591750 |
| 8 | 574 | 19719 | 525600 | 4875786 | 42730578 | 223502230 | 1146718720 | 4312651995 | 16085261781 |
| 9 | 1681 | 97969 | 4196416 | 53597041 | 640749609 | 4245955921 | 27519010816 | 125061956881 | 562969695625 |
| 10 | 5002 | 489219 | 33564800 | 589530846 | 9611072577 | 80672576050 | 660454273024 | 3626791798575 | 19703925162500 |

Table 1: Initial values for the numebr $f(m, n)$ of nonisomorphic irregular fences builon $[1, m] \times[1, n]$

## 4. Lobsters with an $\alpha$-labeling

A lobster is a tree with the property that the removal of all its leaves results in a caterpillar, and a caterpillar is a tree with the property that the removal of all its leaves results in a path. We refer to this path as the central path of the lobster. An alternative definition was given in [13]. Let $P$ be any of the longest paths in a tree $T ; T$ is called a $k$-distance tree if every vertex is at distance at most $k$ from $P$. Thus, paths are 0 -distance trees, caterpillars are 1 -distance trees, and lobsters are 2-distance trees.

It was conjectured by Bermond [14] that all lobsters are graceful. Several families of graceful lobsters are known. Using the construction of Stanton and Zarnke [15] it is possible to obtain a graceful labeling of any lobster constructed by attaching, to every vertex of a path, a leaf of the star $K_{1, n}$. Burzio and Ferrarese [16] proved that any tree obtained from a graceful tree by replacing each edge with a path of fixed length is graceful. Thus, if the starting tree is a caterpillar and every edge is replaced with a path of length 2 , the resulting graph is a lobster. This is one of the strongest results in this area, the weakest part is that the distance between any two leaves, at distance two, is always even. This problem is solved in the work of Wang et al. [17], as well as in the series of articles of Mishra and Panagrahi [18], [19], [20], and [21]. In all these papers, the lobsters considered share the property that all the vertices in the central path have degree larger than two and the subtrees attached to them must satisfy some structural conditions. Morgan [13] proved that all lobsters with a perfect matching are graceful. In a similar line, Krop [22] showed the same for lobsters with an almost perfect matching.

In this section we explore lobsters that are path-like trees and how to use the $\alpha$-labeling, given in Section 2, to produce new $\alpha$-labeled lobsters.
Suppose that the path $P_{5 m}$ has been labeled using the labeling in Lemma 2.1, and embedded in the grid $[1, m] \times[1,5]$, as we did in Section 2 . Thus, every column in this embedding is a copy of $P_{5}$; moreover, the labeling of the $i$ th copy of $P_{5}$ is a $d_{i}$-graceful labeling shifted $c_{i}$ units, where $d_{i}=n(m-i)+1$ and

$$
c_{i}= \begin{cases}\frac{n(i-1)}{2} & \text { if } i \text { is odd } \\ \frac{n(i-1)+1}{2} & \text { if } i \text { is even }\end{cases}
$$

We claim that when every copy of $P_{5}$ is replaced by a copy of any caterpillar of diameter four, the result still holds; that is, we can concatenate the central vertices of these caterpillars to obtain a lobster with an $\alpha$-labeling. In Figure 4.1 we show the labeling scheme given by Rosa [1] to get an $\alpha$-labeling of a caterpillar of size $n-1$.
Let $G$ be a caterpillar of diameter 4 and order $n$. If all the leaves of $G$ are deleted, we get the path $P_{3}$; thus, we can use the notation $C\left(n_{1}, n_{2}, n_{3}\right)$ to denote the caterpillar of order $n=n_{1}+n_{2}+n_{3}+3$, obtained from $P_{3}$ by attaching $n_{i}$ pendant vertices to the vertex $v_{i}$ of $P_{3}$. In Figure 4.2 we show an $\alpha$-labeling $f$ of $C\left(n_{1}, n_{2}, n_{3}\right)$ together with the reverse of its complementary labeling.


Figure 4.1: $\alpha$-labeling scheme of a caterpillar of size $n-1$


Figure 4.2: $\alpha$-labelings of $C\left(n_{1}, n_{2}, n_{3}\right)$

Lemma 4.1. The lobster $L$, obtained by connecting with an edge the central vertices of two copies of the caterpillar $C\left(n_{1}, n_{2}, n_{3}\right)$, is an $\alpha$-tree.

Proof. The caterpillar $C\left(n_{1}, n_{2}, n_{3}\right)$ has size $n_{1}, n_{2}, n_{3}+2$; the $\alpha$-labeling $f$ of it has boundary value $\lambda=n_{1}+n_{3}$. Then, we label the first copy of this caterpillar using the labeling $f$, which is transformed into a $(n+1)$-graceful labeling. In this way, its central vertex has label $n_{1}$. The second copy of the caterpillar is originally labeled using $\bar{f}_{r}$, this labeling is shifted $n_{1}+n_{3}+1$ units, thus there is no repetition of labels between both copies. The new label of the central vertex of the second copy is $\left(n_{1}+n_{2}+2\right)+\left(n_{1}+n_{3}+1\right)=n+n_{1}$. Hence, if we connect with an edge the central vertices, this edge will have weight $n$. Therefore, the lobster $L$ is an $\alpha$-tree.

This process can be applied to any number of copies of $C\left(n_{1}, n_{2}, n_{3}\right)$, in the same way that it was applied to any number of copies of $P_{n}$ in Section 2 . Thus, we get the following theorem.

Theorem 4.2. For each $1 \leq i \leq k$, let $G_{i}$ be a copy of the caterpillar $C\left(n_{1}, n_{2}, n_{3}\right)$. If for every $1 \leq i \leq k-1$, the central vertex of $G_{i}$ is connected with an edge to the central vertex of $G_{i+1}$, then the resulting graph is a lobster that admits an $\alpha$-labeling.

In Figure 4.3 we show an example of this construction using the caterpillar $C(2,4,3)$ four times. We must observe that the lobsters obtained using these caterpillars do not have a perfect (or almost perfect) matching.


Figure 4.3: $\alpha$-labeling of a lobster in $\mathscr{G}$

For each $1 \leq i \leq k$, let $G_{i}$ be a copy of the caterpillar $C\left(n_{1}, n_{2}, n_{3}\right)$. The family $\mathscr{G}_{k}$ consists of all lobsters formed connecting with an edge the central vertices of $G_{i}$ and $G_{i+1}$ where $1 \leq i \leq k-1$. Thus we can say that all members of $\mathscr{G}_{k}$ are $\alpha$-trees. Furthermore, for any $G \in \mathscr{G}_{k}$, the $\alpha$-labeling of $G$, obtained using Theorem 4.2, assigns the label 0 to a leaf of $G$ and the label $\lambda$ (when $k$ is odd) or $\lambda+1$ (when $k$ is even) to another leaf, and the distance between these leaves is $k+3$, that is, the diameter of $G$. In [23] we proved that if $B_{1}, B_{2}, \ldots, B_{k}$ is a collection of $\alpha$-labeled blocks, with boundary value $\lambda_{i}$, then the graph obtained amalgamating the vertex labeled 0 in $B_{i}$ with the vertex labeled $\lambda_{i-1}$ in $B_{i-1}$, for every $2 \leq i \leq k$, is an $\alpha$-graph. We refer to this process as the ( $0, \lambda$ )-amalgamation. As we showed before, if $G$ is a caterpillar, there exists an $\alpha$-labeling of $G$ that assigns the labels 0 and $\lambda$ (when the diameter is even) or 0 and $\lambda+1$ (when the diameter is odd) on the leaves of a path of maximum length in $G$. These two properties allow us to prove the following theorem.

Theorem 4.3. Let $G_{1}, G_{2}, \ldots, G_{t}$ be a collection of $\alpha$-graphs, such that $G_{i} \in \mathscr{G}_{k_{i}}$ or $G_{i}$ is a caterpillar. Then, the lobster L, obtained via $(0-\lambda)$-amalgamation of these graphs, is an $\alpha$-tree.

Proof. Suppose that $f_{i}$ is an $\alpha$-labeling of $G_{i}$ with boundary value $\lambda_{i}$. If $G_{i}$ is a caterpillar, we assume that $f_{i}$ is the labeling $f$ in Figure 4.1 . If $G_{i}$ is a lobster in $\mathscr{G}_{k_{i}}$, we assume that $f_{i}$ is the labeling obtained in Theorem 4.2. In both cases, the vertex of $G_{i}$ labeled 0 belongs to a path of maximum length in $G_{i}$. If the vertex of $G_{i}$ labeled $\lambda_{i}$ is on a leaf, then we can identify the vertex labeled 0 in $G_{i+1}$ with the vertex labeled $\lambda_{i}$ in $G_{i}$. The $\alpha$-labeling of the new graph, denoted by $\Gamma_{i+1}$, is obtained by shifting $\lambda_{i}$ units the labeling $f_{i+1}$ and transforming $f_{i}$ into a $d_{i}$-graceful labeling where $d_{i}-1$ is the size of $G_{i+1}$. If the boundary value of this labeling of $\Gamma_{i+1}$ is on a leaf, we concatenate $\Gamma_{i+1}$ with $G_{i+2}$, to obtain an $\alpha$-graph $\Gamma_{i+2}$, and so on until all the amalgamations are done. If the boundary value of this labeling of $\Gamma_{i+1}$ is not on a leaf, then we use the complementary labeling, which puts its boundary value on a leaf, and connect $\Gamma_{i+1}$ with $G_{i+2}$, and continue in this way until all the amalgamations are done. Given the position of the vertices labeled 0 and $\lambda_{i}$, the final graph is a lobster with an $\alpha$-labeling.

In Figure 4.4 we show an example of this construction where $G_{1} \in \mathscr{G}_{2}, G_{2}$ is a caterpillar of size 10 , and $G_{3} \in \mathscr{G}_{3}$.


Figure 4.4: $\alpha$-labeling of a lobster

## 5. Conclusions

There is a wide variety of fences, we explored here one of these varieties where two consecutive copies of $P_{n}$ are connected by one or two links, if two links are used, the distance between them is odd. These constraints can be modified to explore the existence of $\alpha$-labelings of general fences, where the number of links is not restricted to 1 or 2 . We think that all fences admit an $\alpha$-labeling, except when the fence is isomorphic to the cycle $C_{n}$ with $n \equiv 2(\bmod 4)$, that is not a graceful graph.

The construction of $\alpha$-lobsters presented in Theorem 4.3 can be use in a more general case, where a lobster could be decomposed into sublobsters, each of them with an $\alpha$-labeling that assigns the labels 0 and $\lambda$ to leaves $u$ and $v$ such that the distance between them equals the diameter of the sublobster. We think that this technique should be explored with more details in future works.

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