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by

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On Equalizers And Coequalizers In Comma Categories

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Maclane [3] has discussed equalizers and coequalizers in comma categories $(\mathcal{C} \downarrow \mathcal{A})$ and $(\mathcal{T} \downarrow \mathcal{A})$; we study their creation, reflection and preservation by the projection functors in the comma categories $(\mathcal{T} \downarrow \mathcal{S})$ and $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$.

1. PRELIMINARIES

Definition 1.1.: Given categories and functors

$$\mathcal{B} \xrightarrow{T} \mathcal{C} \xleftarrow{S} \mathcal{D}$$

the (general) comma category $(\mathcal{T} \downarrow \mathcal{S})$ has as objects all triples (B, D, f) with $B \in |\mathcal{B}|$, $D \in |\mathcal{D}|$ and $f: TB \rightarrow SD$; and as morphisms $(B, D, f) \rightarrow (B', D', f')$ all pairs (u, v) of morphisms such that $f' \circ T u = S v \circ f$.
Diagrammatically:

$$\text{Objects } (B, D, f) : \begin{array}{c} TB \\ \downarrow f \\ SD \end{array}; \text{ morphisms } (u, v) : \begin{array}{ccc} TB & \xrightarrow{T(u)} & TB' \\ \downarrow f & & \downarrow f' \\ SD & \xrightarrow{S(v)} & SD' \end{array}$$

The composition $(u', v') \circ (u, v)$ is $(u' \circ u, v' \circ v)$, when defined.

Definition 1.2: Let \mathcal{C} be an arbitrary category and assume that Cat denotes the category of small categories; then the category

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(Cat $\downarrow \cdot \mathcal{C}$) known as super-comma category in the sense of MacLane [3], has objects all pairs (J, F) where J is a small category and $F: J \rightarrow \mathcal{C}$ a functor; and morphisms $(W, \beta): (J, F) \rightarrow (J', F')$ are those pairs consisting of a functor $W: J \rightarrow J'$ and a natural transformation $\beta: F'W \rightarrow F$. Diagrammatically we have

$$\text{Objects } (J, F) : \begin{array}{c} J \\ \downarrow F \\ \mathcal{C} \end{array} ; \text{ morphisms } (W, \beta) : \begin{array}{ccc} J & \xrightarrow{W} & J' \\ & \searrow F & \swarrow F' \\ & & \mathcal{C} \end{array}$$

β

The composition $(W', \beta') \circ (W, \beta)$ is given by $(W'W, \beta' \beta W)$.

The category is also known as 'large diagram category' in the sense of Pareigis [4]. In fact, the concept of this category was introduced by Eilenberg and MacLane [1] pp. 277-280 and later on, was generalized by D.M. Kan [2].

2. COMMA CATEGORY $(T \downarrow S)$

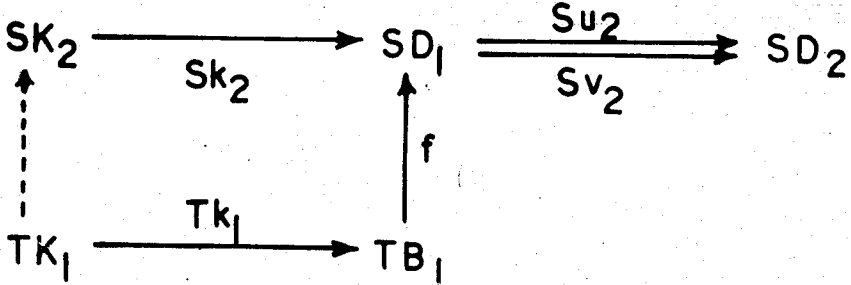
Let us define a functor $Q: (T \downarrow S) \rightarrow \mathcal{B} \times \mathcal{D}$ as: Q assigns to each object (B, D, f) in $(T \downarrow S)$ the object (B, D) in $\mathcal{B} \times \mathcal{D}$ and to each morphism $(u, v): (B, D, f) \rightarrow (B', D', f')$, the morphism $(u, v): (B, D) \rightarrow (B', D')$.

We now prove that

Theorem 2.1: $Q: (T \downarrow S) \rightarrow \mathcal{B} \times \mathcal{D}$ creates equalizers if \mathcal{D} is a category with equalizers and S preserves equalizers.

Proof: Let $(B_1, D_1, f) \xrightarrow{(u_1, u_2)} (B_2, D_2, g) \rightrightarrows (B_1, D_1, f)$ be two morphisms in $(T \downarrow S)$ and let (K_1, K_2) with (k_1, k_2) be the equalizer of (B_1, D_1)

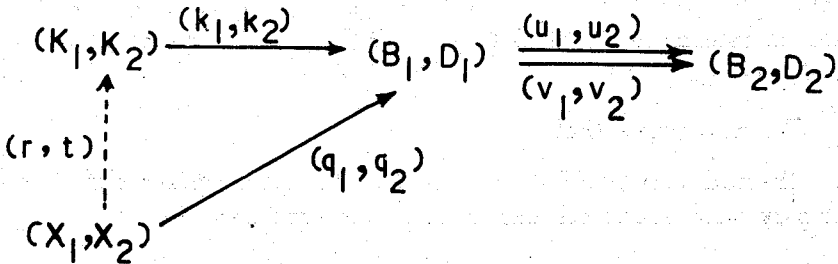
$\xrightarrow[(v_1, v_2)]{(u_1, u_2)} (B_2, D_2)$ in $\mathcal{B} \times \mathcal{D}$. If we consider



then, as S preserves equalizers $(SK_2; Sk_2)$ is the equalizer of Su_2, Sv_2 and moreover $Su_2 f T k_1 = Sv_2 f T k_1$; therefore by definition of equalizer, there exists a unique morphism, say, $k: TK_1 \rightarrow SK_2$ making the above diagram commutative. It yields that (K_1, K_2, k) is an object of $(T \downarrow S)$ and (k_1, k_2) with $(u_1, u_2)(k_1, k_2) = (v_1, v_2)(k_1, k_2)$, is a morphism in $(T \downarrow S)$. We claim that (K_1, K_2, k) with (k_1, k_2) is the

equalizer of $(B_1, D_1, f) \xrightarrow[(v_1, v_2)]{(u_1, u_2)} (B_2, D_2, g)$. Consider (X_1, X_2, h) and

$(q_1, q_2): (X_1, X_2, h) \rightarrow (B_1, D_1, f)$ with $(u_1, u_2)(q_1, q_2) = (v_1, v_2)(q_1, q_2)$. Then if we consider



there is a unique morphism, say, $(r, t): (X_1, X_2) \rightarrow (K_1, K_2)$ making the above diagram commutative. Moreover, we have

$$\text{Sk}_2\text{kTr} = \text{fTk}_1\text{Tr} = \text{f Tq}_1 = \text{Sq}_2\text{h} = \text{Sk}_2\text{Sth}$$

thus it follows that $\text{kTr} = \text{Sth}$ and therefore $(r,t) \in (\text{T} \downarrow \text{S})$. Hence Q creates equalizers.

Dually, we can prove the following:

Theorem 2.2: $\text{Q}: (\text{T} \downarrow \text{S}) \longrightarrow \mathcal{B} \times \mathcal{D}$ creates coequalizers if \mathcal{B} is a category with coequalizers and T preserves coequalizers.

The following corollaries can be deduced from these theorems.

Corollary 2.3: $\text{Q}: (\text{T} \downarrow \text{S}) \longrightarrow \mathcal{B} \times \mathcal{D}$ reflects equalizers and coequalizers.

Corollary 2.4: $\text{Q}: (\text{T} \downarrow \text{S}) \longrightarrow \mathcal{B} \times \mathcal{D}$ preserves equalizers if Q creates equalizers and S preserves those equalizers; Dually, for coequalizers.

3. COMMA CATEGORY $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$

Define a functor $\text{Q}: (\text{Cat} \cdot \downarrow \cdot \mathcal{C}) \longrightarrow \text{Cat}$ such that Q sends each object (J, F) of $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$ to J of Cat and each morphism $(\text{W}, \beta): (\text{J}, \text{F}) \longrightarrow (\text{J}', \text{F}')$ to $\text{W}: \text{J} \longrightarrow \text{J}'$.

Now we show that.

Theorem 3.1: $\text{Q}: (\text{Cat} \cdot \downarrow \cdot \mathcal{C}) \longrightarrow \text{Cat}$ creates equalizers if \mathcal{C} is a category with coequalizers.

Proof: Let $(\text{J}_1, \text{F}_1) \xrightarrow{(\text{W}_1, \beta_1)} (\text{J}_2, \text{F}_2) \xrightarrow{(\text{W}_2, \beta_2)}$ be two morphisms in

$(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$ and $(\text{J}; \text{W})$ be the equalizer of $\text{Q}(\text{W}_1, \beta_1)$, $\text{Q}(\text{W}_2, \beta_2)$ i.e. of W_1 and W_2 in Cat ; then $\text{W}_1\text{W} = \text{W}_2\text{W}$. Since $\beta_1: \text{F}_2\text{W}_1 \longrightarrow \text{F}_1$ and $\beta_2: \text{F}_2\text{W}_2 \longrightarrow \text{F}_1$ are natural transformations, it follows that so are $\beta_1\text{W}: \text{F}_2\text{W}_1\text{W} \longrightarrow \text{F}_1\text{W}$, $\beta_2\text{W}: \text{F}_2\text{W}_2\text{W} \longrightarrow \text{F}_1\text{W}$. As \mathcal{C} is a category with coequalizers, $[\text{J}, \mathcal{C}]$ (the category of functors from J to \mathcal{C}) is also a category with coequalizers. Let $(\text{K}; \beta)$ be the coequalizer of $\beta_1\text{W}$, $\beta_2\text{W}$, i.e. $\beta \beta_1\text{W} = \beta \beta_2\text{W}$; so that (J, K) is an object in $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$ and $(\text{W}, \beta): (\text{J}, \text{K}) \longrightarrow (\text{J}_1, \text{F}_1)$ is a morphism in $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$ with $(\text{W}_1, \beta_1) \circ (\text{W}, \beta) = (\text{W}_2, \beta_2) \circ (\text{W}, \beta)$. Consider (J_3, F_3) in $(\text{Cat} \cdot \downarrow \cdot \mathcal{C})$ and a morphism $(\text{W}_3, \beta_3): (\text{J}_3, \text{F}_3) \longrightarrow (\text{J}_1, \text{F}_1)$ such that $(\text{W}_1, \beta_1) \circ (\text{W}_3, \beta_3) = (\text{W}_2, \beta_2) \circ (\text{W}_3, \beta_3)$; i.e.

$(W_1 \times W_3, \beta_3 \times \beta_1 \times W_3) = (W_2 \times W_3, \beta_3 \times \beta_2 \times W_3)$. Hence by definition of equalizer there is a unique morphism (functor) $X: J_3 \rightarrow J$ in Cat such that $WX = W_3$. Replace W_3 by WX and for any $j \in J_3$, consider the following diagram

$$\begin{array}{ccccc}
 F_2 W_1 W(X(J)) & \xrightarrow[\beta_2 W(X(J))]{\beta_1 W(X(J))} & F_1 W(X(J)) & \xrightarrow{\beta(X(J))} & K(X(J)) \\
 & & & \searrow \beta_3(j) & \downarrow \alpha(j) \\
 & & & & F_3(j)
 \end{array}$$

Then there exists a unique natural transformation $\alpha: KX \rightarrow F_3$ such that $\alpha \beta X = \beta_3$. We have thus obtained a unique morphism $(X, \alpha): (J_3, F_3) \rightarrow (J, K)$ such that $(W, \beta) \circ (X, \alpha) = (W_3, \beta_3)$

$$\begin{array}{ccccc}
 (J, K) & \xrightarrow{(W, \beta)} & (J_1, F_1) & \xrightarrow[\beta_2 W_2]{\beta_1 W_1} & (J_2, F_2) \\
 \uparrow (X, \alpha) & & \nearrow (W_3, \beta_3) & & \\
 (J_3, F_3) & & & &
 \end{array}$$

The following can be proved as Corollaries to the above theorem.

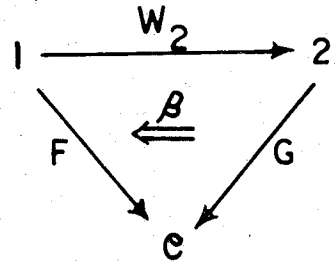
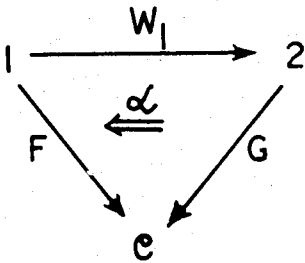
Corollary 3.2: $Q: (\text{Cat} \downarrow \cdot \downarrow \cdot \mathcal{C}) \rightarrow \text{Cat}$ reflects equalizers if \mathcal{C} is a category with coequalizers.

Corollary 3.3: $Q: (\text{Cat} \downarrow \cdot \downarrow \cdot \mathcal{C}) \rightarrow \text{Cat}$ preserves equalizers if \mathcal{C} is a category with coequalizers.

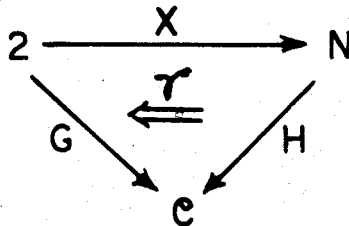
However, the corresponding results for coequalizers do not hold. We give an example in support of our statement.

Example 3.4: Suppose that \mathcal{C} is the category $C_1 \xrightleftharpoons[p]{p} C_2$. Let

1 be the unit category, and 2 the category $A \xrightarrow{f} B$. Let $W_1, W_2: 1 \rightarrow 2$ be the functors sending 1 to A and B respectively; their coequalizer in Cat is $X: 2 \rightarrow \mathcal{N}$, where \mathcal{N} has one object $*$ and its morphisms are the free abelian monoid on t , say, and $X(f) = t$. Let $F: 1 \rightarrow \mathcal{C}$ send 1 to C_2 ; and $G: 2 \rightarrow \mathcal{C}$ be the constant functor at C_1 ; and let natural transformations α and β



be given by $\alpha 1 = p$, $\beta 1 = q$. If \mathcal{C} created coequalizers, the coequalizer of (W_1, α) and (W_2, β) would have the form



This is impossible, for $\alpha \cdot \gamma W_1$ and $\beta \cdot \gamma W_2$ are different.

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