SIMULATION MODELS OF CALL ADMISSION CONTROL SCHEMES USING GPSS

Vassilya ABDULOVA¹, Işık AYBAY²

¹Istanbul Aydin University, Istanbul, TURKEY ²Eastern Mediterranean University, TRNC via Mersin 10, TURKEY

vassilyaabdulova@aydin.edu.tr, isik.aybay@emu.edu.tr

Abstract: In cellular wireless networks, a variety of channel allocation schemes have been developed for achieving high capacity with minimal interference. The choice of channel allocation scheme impacts the performance of the system, particularly as how calls are managed when a mobile user is handed off from one cell to another. Call Admission Control schemes take into account the effect of handoffs in the performance of the system, particularly and call dropping probability. In this study, we present simulation models and programs of some popular Call Admission Control schemes using GPSS simulation tool.

Keywords: GPSS, Call Admission Control, Modelling using simulation.

1. Introduction

Mobile communication systems normally contain three key constituent elements: *mobile station* (MS), *base station* (BS), and *mobile switching center* (MSC). The MS may be any handheld device, such as a mobile phone, portable computer, personal digital assistant (PDA), car communication systems, notebook, or any other device capable of communicating via omni-directional radio waves within a given frequency band using a prescribed communication protocol [1-17]. The geographical area covered by a mobile communication network is divided into several regions known as cells, as shown in figure 1. Each cell has either one or more antennae at its center that are wire-connected to a BS. The BS controls the communications between all MSs inside a particular cell. To facilitate communications beyond the coverage of a single cell, all BSs within a large geographical area are controlled by an MSC, which stores information relevant to all MSs under its control and acts as a gateway to external networks such as the

Internet, public switched telephone network (PSTN), or any other MSC. BSs are normally connected to an MSC by high-speed optical fiber transmission lines [1-17].



Figure 1: Architecture of a cellular network.

The radius of a cell is primarily governed by the power level of the signal. When two cells overlap, an imaginary handoff boundary is assumed along the common chord. In the ideal case, when the radii of all neighboring cells are assumed to be the same, the shape of each cell, formed by the imaginary handoff boundary lines, is a perfect hexagon. In reality, the radius of a cell varies in different directions due to the terrain, co-channel interference effects, and noise. Moreover, depending on location and capacity requirements, different signal power levels may be used neighboring cells, leading to a in heterogeneous cell configuration [1-17].

When a mobile user moves from one cell to another, a handoff occurs. With the active participation of the mobile hosts and the MSC, the base stations are instrumental in initiating and finalizing handoffs. One of the key elements in providing OoS guarantees is to provide a faultless handoff. When there is no channel available for handoff connection, the connection will be dropped. Dropping handoff connection is generally a considered to have a more negative impact from users' perception than rejecting a newly requested connection [1-17].

The mobile hosts in a cell communicate directly with the corresponding BS that has the responsibility of handling all demands for service originating in the cell. In particular, the BS is in charge of negotiating QoS parameters, of performing admission control, and of reserving resources for ongoing connections. This could mean denying access to new connections in order to provide an acceptable level of service to active connections. A new connection that is denied access into the network is said to be blocked. The call blocking probability (CBP) is an important QoS parameter in wireless networks [1-17].

The call dropping probability (CDP) is the possibility that the network will deny resources to a connection after it is in progress. A connection may be dropped during a handoff, when the host moves from a cell into a new cell that is too congested to support more traffic [1-17].

A good scheme must minimize handoff dropping probability because from mobile user's perspective dropping an ongoing connection is less desirable than blocking a new connection. Moreover, the blocking probability of new calls should be reduced and channel utilization maximized [1-17].

Call Admission Control refers to the task of deciding whether or not a certain connection request will be admitted into, and supported by, the network. Prioritized call admission control schemes restrict the number of new calls accepted to decrease the probability of handoff call failure [1-17].

Simulation Models

A computer simulation is a computation that models some real or imagined system over time. Computer simulation makes it possible to analyze systems that would be expensive, dangerous or even impossible to construct prototypes for. The goal of simulation is to evaluate performance characteristics of the system. The behavior of a system as it evolves over time is studied by developing a simulation model. This model usually takes the form of a set of assumptions concerning the operation of the system. Simulation can be used in the design of a new system to check the design variants before implementation. It can also be used in the analysis of an existing system to investigate different modes of operation outside the real system.

In this study, a GPSS World simulation tool was used to simulate some of the most popular Call Admission Control schemes. GPSS World is based on the seminal language of computer simulation, GPSS, which stands for General Purpose Simulation System [18]. This language was developed primarily by Geoffrey Gordon at IBM around 1960, and has contributed important concepts to every commercial discrete event Computer Simulation Language developed ever since. GPSS World is a direct descendent of GPSS/PC, an early implementation of GPSS for personal computers. Since its introduction in 1984, GPSS/PC and its successors have saved thousands of users millions of dollars. GPSS World is the worthy descendent of these early simulation environments.

In the following sections, simulation models and simulation programs for some popular Call Admission Control schemes are presented. It is assumed that voice calls require 1 channel.

Fully Shared Scheme (FSS)

In this scheme, handoff and new calls are treated equally. All available channels are shared by handoff and new calls. This scheme does guarantee the required dropping probability for handover calls, but minimizes call dropping probability and has the advantage of efficient utilization. Flowchart of the simulation model is given in figure 2.



Figure 2: Simulation model of FSS.

Simulation program of FSS:

IAINT	EQU 1
MST	EQU 180

TCH STORAGE 100 GENERATE (EXPONEN TIAL (1,0,IAINT) TRANSFER .5,NC,HO

NC	TEST GE X\$ACH,100,NCAC
	SAVEVALUE NCB+,1

TERMINATE 1

NCAC	SAVEVALUE NCA+,1
	SAVEVALUE ACH+,1
	ENTER TCH,1
	ADVANCE
(EXPONENT	IAL(1,0,MST))
	LEAVE TCH,1
	SAVEVALUE ACH-,1
	TERMINATE 1

- HO TEST GE X\$ACH,100,HOAC SAVEVALUE HOB+,1 TERMINATE 1
- HOAC SAVEVALUE ACH+,1 SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 SAVEVALUE ACH-,1 TERMINATE 1

1.1 Fully Shared Scheme with

Handoff Queuing (FSS_HOQ)

This scheme is an extension of FSS. In this scheme, priority is given to handoff attempts by allowing queuing of handoff calls in case all channels are occupied. Queuing of handoff calls will result in a lower call dropping probability compared to FSS. When a channel is released in a cell, it is assigned to the next handoff call attempt waiting in the queue (if any). If more than one handoff call attempt is in the queue, the first-come-first-served queuing discipline is used. If no channel was allocated to the handoff call in the queue, it will be dropped after a specific

period of time, td. So, if time spent by handoff in the queue, th, is greater or equal to td, then handoff call is dropped. A new call is blocked if no channel is available in the cell. Flowchart of the simulation model is given in figure 3.



Figure 3: Simulation model of FSS_HOQ scheme.

Simulation program of FSS_HOQ:

IAINT	EQU 1
MST	EQU 180
HQINT	EQU 20

TCH STORAGE 100
GENERATE
(EXPONENTIAL(1,0,IAINT))
TRANSFER .5.NC.HO

NC	TEST GE X\$ACH,100,NCAC
NCBL	SAVEVALUE NCB+,1
	TERMINATE 1

NCAC SAVEVALUE NCA+,1 SAVEVALUE ACH+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,NCNOQH SAVEVALUE ACHQH+,1

SAVEVALUE QH-,1

TERMINATE 1

NCNOQH	SAVEVALUE	ACH-,1
	TERMINATE	1

HO TEST GE X\$ACH,100,HOAC QUEUE Handoff MARK SAVEVALUE QH+,1

ISCHAV ADVANCE 0.1 TEST GE X\$ACHQH,1,NCHAV CHAV SAVEVALUE ACHQH-,1

ENTER TCH,1 DEPART Handoff ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,QNOQH SAVEVALUE ACHQH+,1

> SAVEVALUE QH-,1 TERMINATE 1

QNOQH SAVEVALUE ACH-,1 TERMINATE 1

TEST GE NCHAV M1,HQINT,ISCHAV **DEPART Handoff** SAVEVALUE QH-,1 SAVEVALUE HOB+,1 HODR **TERMINATE 1** HOAC SAVEVALUE ACH+,1 SAVEVALUE HOA+,1 ENTER TCH.1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,HNOQH SAVEVALUE ACHQH+,1 SAVEVALUE QH-,1 **TERMINATE 1**

HNOQH SAVEVALUE ACH-,1 TERMINATE 1

1.2 Fully Shared Scheme with

Handoff and New Call Queuing

(FSS_HOQ_NCQ)

This scheme is extension an of FSS_HOQ. In this scheme, both handoff and new calls are queued. Queuing of new calls will result in a lower call blocking probability and higher call dropping probability compared to FSS_HOQ. In this scheme handoff calls are treated in same way as in FSS_HOQ. A new call is queued if no channel is available in the cell. When a channel is released in a cell, it is assigned to the next handoff call attempt waiting in the queue. If there is no handoff waiting in the queue, the channel is assigned to the next new call waiting in the queue (if any). If more than one new call is in the queue, the first-come-firstserved queuing discipline is used. If no channel was allocated to the new call in the queue, it will be blocked after a specific period of time, tb. So, if time spent by a new call in the queue, tn, is greater or equal to tb, then new call is blocked. Flowchart of the simulation model is given in figure 4.



Figure 4: Simulation model of FSS_HOQ_NCQ scheme.

Simulation	program of	
FSS_HOQ	_NCQ:	
IAINT	EQU 1	
MST	EQU 180	
HQINT	EQU 20	
NCINT	EQU 20	
	TCH STORAGE 100	
	GENERATE	
(EXPONENTIAL(1,0,IAINT))		
	TRANSFER .5,NC,HO	
NC	TEST GE X\$ACH,100,NCAC	
	QUEUE Newcall	
	MARK	
NCISCHAV	ADVANCE 0.1	
	TEST GE	
X\$ACH,100,NCCHAV		
	TEST GE	
M1,NCINT,N	CISCHAV	
	DEPART Newcall	
NCBLCK	SAVEVALUE NCB+,1	
	TERMINATE 1	
NCCHAV	SAVEVALUE ACH+,1	
	ENTER TCH.1	

DEPART Newcall

ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,NCNOQH SAVEVALUE ACHQH+,1

> SAVEVALUE QH-,1 TERMINATE 1

NCAC SAVEVALUE ACH+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,NCNOQH SAVEVALUE ACHQH+,1

> SAVEVALUE QH-,1 TERMINATE 1

NCNOQH SAVEVALUE ACH-,1 TERMINATE 1

HO TEST GE X\$ACH,100,HOACPT QUEUE Handoff MARK SAVEVALUE QH+,1

HOISCHAV ADVANCE 0.1 TEST GE X\$ACHQH,1,HONOCHAV HOCHAV SAVEVALUE ACHQH-,1

ENTER TCH,1 DEPART Handoff ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,HONOQH SAVEVALUE ACHQH+,1

> SAVEVALUE QH-,1 TERMINATE 1

HONOCHAV TEST GE M1,HQINT,HOISCHAV DEPART Handoff SAVEVALUE QH-,1 HODROP SAVEVALUE HOB+,1 TERMINATE 1

HOACPT SAVEVALUE ACH+,1

SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,HONOQH SAVEVALUE ACHQH+,1

> SAVEVALUE QH-,1 TERMINATE 1

HONOQH SAVEVALUE ACH-,1 TERMINATE 1

1.3 Guard Channel Scheme (GCS)

In this scheme priority is given to handoff attempts by assigning Ch channels exclusively for handoff calls among the C channels in a cell. The remaining C – Ch channels are shared by both new calls and handoff calls. A new call is blocked if the number of available channels in the cell is less than or equal to Ch when the call is originated. Α handoff attempt is unsuccessful if no channel is available in the target cell. Flowchart of the simulation model is given in figure 5. In the flowchart, C represents total number of channels, Ca represents number of already allocated channels and Ch represents number of channels reserved for handoff calls.



Figure 5: Simulation model of GCS.

Simulation program of GCS:

IAINT MST	EQU 1 EQU 180
(EXPONENTI	TCH STORAGE 100 GENERATE (AL(1,0,IAINT)) TRANSFER .5,NC,HO
NC	TEST GE X\$ACH,80,NCAC SAVEVALUE NCB+,1 TERMINATE 1
NCAC	SAVEVALUE NCA+,1 SAVEVALUE ACH+,1 ENTER TCH,1 ADVANCE
(EXPONENTI	IADVANCE IAL(1,0,MST)) LEAVE TCH,1 SAVEVALUE ACH-,1 TERMINATE 1
НО	TEST GE X\$ACH,100,HOAC SAVEVALUE HOB+,1 TERMINATE 1

HOAC SAVEVALUE ACH+,1 SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 SAVEVALUE ACH-,1 TERMINATE 1

1.4 Guard Channel Scheme with

Handoff Queuing (GCS_HOQ)

In this scheme, same channelsharing method is used as in GCS, except that queuing of handoff calls is allowed if necessary. Queuing of handoff calls will result in a lower call dropping probability compared to GCS. If there is no available channel for handoff call in a cell, handoff call is queued. When a channel is released in a cell, it is assigned to the next handoff call attempt waiting in the queue (if any). If more than one handoff call attempt is in the queue, the first-come-first-served queuing discipline is used. If no channel was allocated to the handoff call in the queue, it will be dropped after a specific period of time, td. So, if time spent by handoff in the queue, th, is greater or equal to td, then handoff call is dropped. A new call is handled the same way as in GCS. Flowchart of the simulation model is given in figure 6. In the flowchart, C represents total number of channels, Ca represents number of already allocated

channels and Ch represents number of channels reserved for handoff calls.



Figure 6: Simulation model of GCS_HOQ

scheme.

Simulation program of GCS_HOQ scheme:

IAINT	EQU 1
MST	EQU 180
HQINT	EQU 20

TCH STORAGE 100 GENERATE (EXPONENTIAL(1,0,IAINT)) TRANSFER .5,NC,HO

NC	TEST GE X\$ACH,80,NCAC
NCBL	SAVEVALUE NCB+,1
	TERMINATE 1

NCAC SAVEVALUE NCA+,1 SAVEVALUE ACH+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,NCNOQH SAVEVALUE ACHQH+,1

SAVEVALUE QH-,1 TERMINATE 1

- NCNOQH SAVEVALUE ACH-,1 TERMINATE 1
- HO TEST GE X\$ACH,100,HOAC QUEUE Handoff MARK SAVEVALUE QH+,1
- ISCHAV ADVANCE 0.5 TEST GE X\$ACHQH,1,NCHAV CHAV SAVEVALUE ACHQH-,1
- ENTER TCH,1 DEPART Handoff ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,QNOQH SAVEVALUE ACHQH+,1

SAVEVALUE QH-,1 TERMINATE 1

- QNOQH SAVEVALUE ACH-,1 TERMINATE 1
- NCHAV TEST GE M1,HQINT,ISCHAV DEPART Handoff SAVEVALUE QH-,1 HODR SAVEVALUE HOB+,1 TERMINATE 1
- HOAC SAVEVALUE ACH+,1 SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,HNOQH SAVEVALUE ACHQH+,1

SAVEVALUE QH-,1 TERMINATE 1

HNOQH SAVEVALUE ACH-,1 TERMINATE 1

1.5 Guard Channel Scheme with

New Call Queuing (GCS_NCQ)

In this scheme, same channel-sharing method is used as in GCS, except that queuing of new calls is allowed if necessary. Queuing of new calls will result in a lower call blocking probability compared to GCS. If there is no available channel for the new call, new call is queued. When a channel is released in a cell, it is assigned to the next new call waiting in the queue (if any). If more than one new call is in the queue, the firstcome-first-served queuing discipline is used. If no channel was allocated to the new call in the queue, it will be blocked after a specific period of time, tb. So, if time spent by new call in the queue, tn, is greater or equal to tb, then new call is blocked. Flowchart of the simulation model is given in figure 7. In the flowchart, C represents total number of channels, Ca represents number of already allocated channels and Ch represents number of channels reserved for handoff calls.





scheme.

Simulation program of GCS_NCQ scheme:

IAINT	EQU 1
MST	EQU 180
NCINT	EQU 20
	TCH STORAGE 100
	GENERATE
(EXPONENT	IAL(1,0,IAINT))
	TRANSFER .5,NC,HO
NC	TEST GE X\$ACH,80,NCAC
	QUEUE Newcall
	MARK
NCISCHAV	ADVANCE 0.1
	TEST GE
X\$ACH,80,N	CCHAV
. , ,	TEST GE
M1,NCINT,N	CISCHAV
	DEPART Newcall
NCBLCK	SAVEVALUE NCB+,1
	TERMINATE 1
NCCHAV	SAVEVALUE ACH+,1
	ENTER TCH,1
	DEPART Newcall
	ADVANCE
(EXPONENT	IAL(1,0,MST))
	LEAVE TCH,1

SAVEVALUE ACH-,1 SAVEVALUE NCA+,1 TERMINATE 1

NCAC SAVEVALUE NCA+,1 SAVEVALUE ACH+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 SAVEVALUE ACH-,1 TERMINATE 1

HO TEST GE X\$ACH,100,HOAC HODROP SAVEVALUE HOB+,1 TERMINATE 1

HOAC SAVEVALUE ACH+,1 SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE (EXPONENTIAL(1,0,MST)) LEAVE TCH,1 SAVEVALUE ACH-,1 TERMINATE 1

1.6 Guard Channel Scheme with

Handoff and New Call Queuing

(GCS_HOQ_NCQ)

This scheme is a combination of GCS_HOQ and GCS_NCQ. In this scheme, same channel-sharing method is used as in GCS, except that queuing of both handoff and new calls is allowed if necessary. When a channel is released in a cell, it is assigned to the next handoff call attempt waiting in the queue. If there is no handoff waiting in the queue, the channel is assigned to the next new call waiting in the queue (if any). The queuing of handoff

calls is performed in the same manner as in GCS_HOQ and queuing of new calls is performed in the same manner as in GCS_NCQ. Flowchart of the simulation model is given in figure 8. In the flowchart, C represents total number of channels, Ca represents number of already allocated channels and C_h represents number of channels reserved for handoff calls.



GCS_HOQ_NCQ scheme.

Simulation	program	of
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GCS_HOQ_NCQ scheme:

IAINT	EQU 1		
MST	EQU 180		
HQINT	EQU 20		
NCINT	EQU 20		
TCH STORAGE 100 GENERATE (EXPONENTIAL(1,0,IAINT)) TRANSFER .5,NC,HO			

NC	TEST GE X\$ACH,80,NCAC QUEUE Newcall MARK		TEST GE X\$QH,1,HONOQH SAVEVALUE ACHQH+,1	
NCISCHAV	ADVANCE 0.1 TEST GE		SAVEVALUE QH-,1 TERMINATE 1	
X\$ACH,80,N	CCHAV TEST GE			
M1,NCINT,NCISCHAV DEPART Newcall		HONOCHAV TEST GE M1,HQINT,HOISCHAV		
NCBLCK	SAVEVALUE NCB+,1 TERMINATE 1	HODROP	DEPART Handoff SAVEVALUE QH-,1 SAVEVALUE HOB+,1	
NCCHAV	SAVEVALUE ACH+,1 ENTER TCH,1 DEPART Newcall	HOACPT	TERMINATE 1 SAVEVALUE ACH+,1	
(EXPONENT	ADVANCE TAL(1,0,MST)) LEAVE TCH,1		SAVEVALUE HOA+,1 ENTER TCH,1 ADVANCE	
	TEST GE X\$QH,1,NCNOQH SAVEVALUE ACHQH+,1 SAVEVALUE OH- 1	(EXPONEN	FIAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,HONOQH SAVEVALUE ACHOH+ 1	
NCAC	TERMINATE 1 SAVEVALUE ACH+,1 ENTER TCH,1		SAVEVALUE QH-,1 TERMINATE 1	
(EXPONENT	ADVANCE IAL(1,0,MST)) LEAVE TCH,1 TEST GE X\$QH,1,NCNOQH	HONOQH	SAVEVALUE ACH-,1 TERMINATE 1	
	SAVEVALUE ACHQH+,I	2. Co	2. Conclusion	
	SAVEVALUE QH-,1 TERMINATE 1	In this stud	y, GPSS was used in order to	
NCNOQH	SAVEVALUE ACH-,1 TERMINATE 1	model some of the popular Call Admission Control schemes for cellular		
HO X\$ACH,100,I	TEST GE HOACPT QUEUE Handoff MARK SAVEVALUE QH+,1	wireless networks. The following schemes dealing were modelled: Fully Shared Scheme (FSS), Fully Shared Scheme with Handoff Queuing (FSS HQQ) Fully		
HOISCHAV	ADVANCE 0.1 TEST GE	Shared Scheme with Handoff and New		
X\$ACHQH,1 HOCHAV	,HONOCHAV SAVEVALUE ACHQH-,1	Call Queui Channel Sc	ng (FSS_HOQ_NCQ), Guard channel	
	ENTER TCH,1 DEPART Handoff ADVANCE	Scheme	with Handoff Queuing)), Guard Channel Scheme	
(EXPONENT	IAL(1,0,MST)) LEAVE TCH,1	(365_110)		

with New Call Queuing (GCS_NCQ) and Guard Channel Scheme with Handoff and New Call Queuing (GCS_HOQ_NCQ).

This paper will be useful for academia interested in modelling of CAC schemes using GPSS.

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