# Solving Multidimensional Knapsack Problem with Bayesian Multiploid Genetic Algorithm 

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#### Abstract

Solving optimization problems is still a big challenge in the area of optimization algorithms. Many proposed algorithms in the literature don't consider the relations between the variables of the nature of the problem. However, a recently published algorithm, called "Bayesian Multiploid Genetic Algorithm" exploits the relations between the variables and then solves the given problem. It also uses more than one genotype unlike the simple Genetic Algorithm (GA) and it acts like an implicit memory in order to remember the old but good solutions. In this work, the well-known Multidimensional Knapsack Problem (MKP) is solved by the Bayesian Multiploid Genetic Algorithm. And the results show that exploiting relations between the variables gets a huge advantage in solving the given problem.


## 1. Introduction

Constraint Optimization Problems (COP) are still a big challenge in the area of computer science. One of them is the Multidimensional Knapsack Problem (MKP). The MKP is an extended version of the standard 0-1 Knapsack Problem (KP). While standard KP has only one resource, the MKP can has more than one resource. MKP is actually can be considered a real-world problem. Many real-world problems can be solved by MKP, such as cutting stock[1], loading problems [2], resource allocation for distributed computing[3], project selection[4], etc. MKP is still a current benchmark problem that continues to be solved with different approaches in recent years[5]-[8].

In the past ten years, many metaheuristic (MH) techniques are applied to solve different NP-Hard global optimization problems. Some of them are the following: Estimation of Distribution Algorithms (EDA)[9], Artificial Bee Colony (ABC)[10], Harmony Search (HS)[11], Ant Colony Optimization
(ACO)[12], Whale Optimization Algorithm (WOA)[13], Bat Algorithm (BA)[14], etc. However, all of them mentioned before don't consider the interactions between the problem's variables. In an ecosystem, all the objects have connected some way and they affect each other.

For this reason, it is important to exploit interactions between the variables and then use them to solve the problem effectively.

In this work, the Bayesian Multiploid Genetic Algorithm (BMGA)[15] is used, which has both an implicit memory scheme (to remember old solutions) and a Bayesian Network (to exploit interactions between variables) in order to solve the well-known Multidimensional Knapsack Problem (MKP) considering as a real-world benchmark problem.

To evaluate the BMGA's performance, six algorithms are selected as comparative algorithms. Note that, the results for the six algorithms listed below are already taken from the [16] and compared

[^0]to the BMGA's results.

- Moth Search Algorithm (MS)[17, s.]: A new optimization algorithm that is inspired by the Levy flights and the phototaxis of the months. In this method, the fittest individual is considered the light source. The moths close to the fittest one fly in the form of Levy flights. On the other hand, because of the phototaxis, the moths far from the fittest one fly to the fittest one with big steps. These two different behaviors are the exploration and the exploitation of any other optimization method.
- Self-Learning MS (SLMS)[16]: In regular MS, each individual update their positions according to the fittest one. But it may cause it to fall into the local optimum. Authors, introduce a new selflearning strategy to enable individuals to update their positions not only according to the fittest one but also the closer individuals which have fitter than themselves.
- Modified multi-verse optimization algorithm (MMVO)[18]: In this algorithm, authors are inspired by the popular multi-verse theory which is based on three concepts: the wormhole, the black hole, and the white hole. In the algorithm, each universe is considered a solution candidate. If an individual gets close to the white hole, means it is getting better fitness values, and vice-versa (closing to the black hole), it is getting worse fitness values. And the wormholes are used as a diversification operator in the algorithm to maintain the diversity in the population.
- Binary Gravitational Search Algorithm (BGSA)[19]: This is the binary version of GSA which is based on Newton's laws of gravity and motion: The gravitational force affects the objects and makes them attract each other. In the algorithm, each object is considered an agent, and each one of them has its mass. And their masses define their fitness values. As time passes, objects are attracted by the fitter masses which leads them to get a better fitness value.
- Binary Hybrid Topology Particle Swarm Optimization (BHTPSO)[19]: In the simple $\operatorname{PSO}[20]$, the particles may fall into local optimum if their velocities are zero. Because zero velocity means the particle's fitness value is good and shouldn't be changed. To overcome this, a small value is added to their velocities for

```
Algorithm 1 Pseudocode of BMGA
\(g \leftarrow\) number of genotypes
pop \(\leftarrow \operatorname{init}(p o p S i z e, g)\)
randomly generate a probVec[]
while a termination condition not met do
    genotype2phenotype(pop,probVec[])
    best \(\leftarrow\) evaluate (pop)
    \(B N \leftarrow\) constructBN (pop)
    auxPop \(\leftarrow\) sampleBN \((B N)\)
    probVec[] \(\leftarrow\) constructProbVec(auxPop)
    pop \(\leftarrow\) tournamentSelection(pop)
    pop \(\leftarrow\) uniformCrossover (pop)
    pop \(\leftarrow\) bitwiseMutation (pop)
return best
```

acceleration and enables them to escape from the local optimum.

- Binary Hybrid Topology Particle Swarm Optimization Quadratic Interpolation (BHTPSOQI)[19]: This is the form of BHTPSO which is incorporated with a quadratic crossover operator.
This paper continues as follows: In Section 2, the BMGA is explained. In Section 3, MKP and its datasets are explained. In Section 4, the results are shown and finally, Section 5 concludes the paper.


## 2. Bayesian Multiploid Genetic Algorithm

The BMGA is constructed on the simple GA. However, it differs from simple GA in many ways.

First, the individuals in the simple GA have only one chromosome to represent a candidate solution. Both genetic operators and fitness calculations are done over this chromosome for each individual. However, in BMGA, each individual has two different structures: The genotypes and a phenotype. As we know, in nature, all living things have genotypes and phenotypes. Genotypes are inherited from the parents. However, the phenotype is the one that decides a living thing will look like to what.

The number of genotypes in an individual can be more than one. This feature provides an implicit memory scheme to the algorithm. Because all the genetic operators (crossover - mutation) are executed on the genotypes and the fitness calculations are executed on the phenotype of the individual. In this way, genotypes act like a memory. On the other hand, each individual has only one phenotype and its fitness value is calculated using that phenotype. That is, no matter how many genotypes an individual has, only one phenotype determines its fitness value. The structure of an individual is illustrated in Figure 1.

Second, BMGA uses a well-known Bayesian

Optimization Algorithm (BOA)[21] in it, in order to exploit the interactions between the variables. The BOA starts with a randomly generated population and creates a Bayesian Network (BN) by using this population. Since there is no prior information, BOA uses a greedy algorithm to form different BNs and then measure their quality using a special metric. After the algorithm finds the most suitable BN, then samples new individuals using that final BN. This procedure continues until a termination condition is met.

|  | 0 | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| genotype \#1 | 1 | 1 | 0 | 1 | 1 | 0 |
| genotype \#2 | 0 | 0 | 0 | 1 | 0 | 0 |
| genotype $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |
| genotype \#k | 1 | 1 | 0 | 1 | 0 | 0 |
| Prob. Vector | $\mathbf{0 . 5}$ | $\mathbf{0 . 3}$ | $\mathbf{0 . 7}$ | $\mathbf{0 . 1}$ | $\mathbf{0 . 8}$ | $\mathbf{0 . 9}$ |
| phenotype | $\boldsymbol{?}$ | $?$ | 0 | 1 | $?$ | 0 |

Figure 1 Illustration of an individual
In BMGA, the BOA is used in order to form a probability vector to determine the phenotypes of the individuals in the population. This works like this: First, the GA part of the BMGA randomly generates a population. In this generating part, only the genotypes of the individuals are generated. For instance, if an individual has four genotypes, and the size of the population is 100 , then 400 genotypes are generated. Next, a probability vector is randomly formed to determine the phenotypes of the individuals for the first iteration. The probability vector is formed via genotypes of the best $k \%$ individuals of the population by calculating the probability of being 1 for each gene using Equation 1. After the probability of being 1 is calculated the Equation 1, for each gene in the phenotype, a random real number is generated between 0 and 1 . This generated value is then compared to the corresponding value in the probability vector to determine the value of the phenotype. Assuming that number of genotypes is $g$, this comparison is held by using Equation 2 and Equation 3.

$$
\begin{align*}
& \text { ProbVector }_{i}=\frac{\text { number of ones for the } i^{\text {th }} \text { gene }}{\text { size of the given population }}  \tag{1}\\
& \text { val }=\left(\text { rand }()<\text { ProbVector }_{i}\right) ? 1: 0 \tag{2}
\end{align*}
$$

$$
\text { pheno }_{i}=\left\{\begin{array}{lr}
0, & \sum^{g} \text { geno }_{i}=0  \tag{3}\\
1, & \sum^{g} \text { geno }_{i}=g \\
\text { val, } & \text { other }
\end{array}\right.
$$

Once the phenotypes are determined, the fitness values of all individuals are calculated. After this first iteration, for the latter iterations, a new BN is constructed by using the same rule, and then a second population is sampled via BN. This second population is named auxiliary population, in short, "aux-pop". The aux-pop has only one chromosome like in the simple GA. For the latter iterations, the probability vector is formed by the aux-pop, again, using Equation 1.

After explaining the key points of the BMGA, now we can explain it in general. The general pseudocode of the BMGA is given in Algorithm 1. After the BOA parts of the BMGA are executed, next, the GA part of the BMGA starts to run.

The standard tournament selection (size of $n$ ) is applied in the BMGA's selection phase: Randomly chosen $n$ solution candidates are compared and the fittest one is passed on to the next generation. This operation is performed until the size of the next generation is satisfied. Without recalculating its fitness value, the elite solution candidate from the last iteration is passed on to the current generation (elitism).

The uniform crossover approach is employed for crossover operation. It begins by randomly generating a mask vector in binary and then selecting two individuals, say $i_{1}$, and $i_{2}$, at each iteration. For each variable j (genes in chromosome), if the mask vector's corresponding value is $1, i_{1}$ 's 1 st genotype's $j^{t h}$ variable is swapped with $i_{2}$ 's 2 nd genotype's $j^{t h}$ variable with a probability of $p_{c}$. This procedure is performed on each pair of genotypes in each solution candidate separately.

The simple bitwise mutation is employed as a mutation tool. The genotypes of each solution candidate are inverted with a probability of pm in the bitwise mutation process.

## 3. Preparing Test Environment

### 3.1. Multidimensional Knapsack Problem

MKP is a well-known benchmark problem for optimization algorithms. The goal of the problem finding a subset of given a number of items that obtain the optimal profit value while satisfying the given constraints. In this problem, there are $n$ items, $m$ resources, and $n \times m$ constraints. The formulation of the problem is given in Equation 4 and Equation 5.

$$
\begin{gather*}
\max \quad \sum_{j=1}^{n} p_{j} \times x_{j}  \tag{4}\\
\text { subject to } \sum_{j=1}^{n} r_{i j} \times x_{j} \leq c_{i}, \quad \forall_{i} \in\{1,2, \ldots, m\} \tag{5}
\end{gather*}
$$

where $n$ is the number of items, $m$ is the number of resources (knapsacks), $x_{j}$ denotes the whether the item $j$ is collected, $p_{j}$ is the profit of the item $j, r_{i j}$ is the consumption of the $j^{\text {th }}$ item at the $i^{\text {th }}$ resource and finally $c_{i}$ is the capacity of the $i^{t h}$ resource.

In order to test BMGA for MKPs, the first five problems are taken from the six different datasets which are provided on the well-known OR-LIB[22]. There are 30 different problems in that dataset and each has a different problem size. All problems have a tightness ratio of 0.25 , however, their number of items and number of resources is changing. Also, they are encoded as "cbX- Y " in the result table, which means "Chu-Beasley, dataset X, problem Y". All the necessary information is given in Table 2.

In this work, all 30 datasets are solved by BMGA. Then, the results are compared to the results obtained from [16]

### 3.2. Experimental Studies

For the experimental tests, first, some parameters of BMGA are set. These settings can be seen in Table 1.

Table 1 Parameter settings

| Parameter | Value |
| :--- | :--- |
| number of generations | 1000 |
| Size of the population | 100 |
| Probability of crossover | 1.0 |
| Probability of mutation | 0.03 |
| Tournament size | 4 |
| number of genotypes | 4 |
| Population rate to form BN | 0.1 |

Note that, each value in Table 1 is obtained by conducting a number of preliminary sensitivity tests[15].

For each test, with the same set of seeds, 30 independent runs were performed. For each run, best-of-generation (BOG) is saved and the overall performance is calculated as shown in Equation 6.

$$
\begin{equation*}
\bar{F}_{B O G}=\frac{1}{G} \sum_{i=1}^{G}\left(\frac{1}{N} \sum_{j=1}^{N} F_{B O G_{i j}}\right) \tag{6}
\end{equation*}
$$

where $G$ is the number of generations, $N$ is the number of runs, $F_{B O G_{i j}}$ is the BOG of the $j^{\text {th }}$ run's $i^{\text {th }}$ generation. Last, $\bar{F}_{B O G}$ is the overall offline performance.

## 4. Results

After setting the parameters in Table 1, the BMGA is tested for 30 different MKP datasets. For each
problem, 30 independent runs are executed. The results can be seen in Table 2. Also, the charts of the first problems of each group can be seen in Figure 2.

In the table, "Optimum" is the best-known optimum value so far provided in [22], "Best" is the best solution found during each test, and "Mean" is the average performance for the particular problem. In the "Prob. size" column, the size of the problem is given. For example, " $5 \times 100-0.25$ " means that the problem has 5 resources, 100 items, and a tightness ratio of 0.25 .

Results show that BMGA is capable to solve MKP, is competitive, and obtains better solutions in terms of \#Mean and \#Best for most of the problem instances because of its Bayesian probability vector. More clearly, both in terms of \#Best and \#Mean, BMGA has got better results in 20 of 30 problems.

Although BMGA performs better than the other algorithms, for the test instances with 500 items, it performs slightly worse than the other test instances. That is because having more items means a bigger BN. And since BOA's BN is constructed with a 1-incoming edge rule, it is getting hard to exploit relations between the variables.

## 5. Conclusion

In this paper, the well-known optimization problem, the MKP, is solved by a recently proposed algorithm, the BMGA.

BMGA combines the powers of EDA and MS. By saying EDA, we mean BOA and by saying MS, we mean GA. While the MS part is responsible for the optimization process, the EDA part is responsible for exploiting relations between variables. Since it is a recently proposed algorithm, it has become mandatory to solve well-known optimization benchmark problems such as MKP.

For testing, BMGA is used to solve the most popular optimization dataset library's MKP instances. Then, the results are compared to the most recent paper. To get a fair comparison, the same number of fitness value calculations were done. The results showed that BMGA outperforms even the latest proposed algorithms.

This work tells us that exploiting the relations between the problem variables is important and useful while solving global optimization problems.

Table 2 Results for the test sets

| Prob. size | Prob. | Optimum | Profit | BMGA | MS | SLMS | MMVO | BGSA | BHTPSO | BHTPSO-QI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5x100-0.25 | cb1-1 | 24381 | Best | 24311 | 24253 | 24231 | 24192 | 24152 | 24169 | 24301 |
|  |  |  | Mean | 24072 | 24004 | 24015 | 24050 | 23835 | 23822 | 23821 |
|  | cb1-2 | 24274 | Best | 24274 | 24258 | 24274 | 24274 | 23986 | 24109 | 23944 |
|  |  |  | Mean | 24225 | 23934 | 24145 | 24274 | 23536 | 23657 | 23688 |
|  | cb1-3 | 23551 | Best | 23247 | 23538 | 23538 | 23538 | 23386 | 23435 | 23418 |
|  |  |  | Mean | 23175 | 23272 | 23440 | 23520 | 23041 | 23072 | 23073 |
|  | cb1-4 | 23534 | Best | 23330 | 23256 | 23330 | 23288 | 23172 | 23253 | 23192 |
|  |  |  | Mean | 23289 | 23024 | 23156 | 23120 | 22863 | 22928 | 22923 |
|  | cb1-5 | 23991 | Best | 23952 | 23845 | 23947 | 23947 | 23755 | 23815 | 23774 |
|  |  |  | Mean | 23901 | 23567 | 23800 | 23900 | 23459 | 23473 | 23527 |
| $5 \times 250-0.25$ | cb2-1 | 59312 | Best | 59203 | 58084 | 59107 | 58473 | 57565 | 57814 | 57800 |
|  |  |  | Mean | 58980 | 57369 | 58736 | 58240 | 56554 | 56874 | 56685 |
|  | cb2-2 | 61472 | Best | 61227 | 60248 | 61280 | 60692 | 60057 | 59982 | 59767 |
|  |  |  | Mean | 61185 | 59386 | 61041 | 60390 | 58613 | 58588 | 58680 |
|  | cb2-3 | 62130 | Best | 61831 | 61212 | 61787 | 61702 | 59936 | 60630 | 60524 |
|  |  |  | Mean | 61684 | 59922 | 61476 | 61330 | 58975 | 59234 | 59186 |
|  | cb2-4 | 59463 | Best | 59167 | 58386 | 59101 | 58441 | 57970 | 57736 | 57884 |
|  |  |  | Mean | 58777 | 57752 | 58787 | 58300 | 56744 | 56773 | 56584 |
|  | cb2-5 | 58951 | Best | 58753 | 57755 | 58485 | 58082 | 56959 | 57378 | 57550 |
|  |  |  | Mean | 58566 | 56929 | 58097 | 58300 | 55961 | 56129 | 56361 |
| 5x500-0.25 | cb3-1 | 120148 | Best | 119992 | 116296 | 119914 | 119978 | 111206 | 114493 | 114438 |
|  |  |  | Mean | 119921 | 115444 | 119625 | 119900 | 108930 | 111017 | 111469 |
|  | cb3-2 | 117879 | Best | 116722 | 113732 | 117362 | 115634 | 108522 | 112821 | 112147 |
|  |  |  | Mean | 115888 | 112257 | 116858 | 115400 | 106631 | 109276 | 109247 |
|  | cb3-3 | 121131 | Best | 117859 | 117666 | 120888 | 119156 | 111271 | 114774 | 116099 |
|  |  |  | Mean | 117083 | 116367 | 120711 | 118900 | 109430 | 112035 | 112001 |
|  | cb3-4 | 120804 | Best | 120501 | 116454 | 120030 | 119124 | 111283 | 115828 | 114327 |
|  |  |  | Mean | 119662 | 115396 | 119644 | 118900 | 109062 | 112200 | 111671 |
|  | cb3-5 | 122319 | Best | 119059 | 117900 | 121907 | 121141 | 112391 | 115889 | 117242 |
|  |  |  | Mean | 118284 | 116767 | 121512 | 120800 | 110564 | 112253 | 113364 |
| 10x100-0.25 | cb4-1 | 23064 | Best | 22917 | 22753 | 22835 | 22805 | 22836 | 22905 | 22876 |
|  |  |  | Mean | $22694$ | $22459$ | 22604 | $22700$ | 22334 | 22425 | 22449 |
|  | cb4-2 | 22801 | Best | 22836 | 22611 | 22650 | 22630 | 22441 | 22573 | 22408 |
|  |  |  | Mean | 22541 | 22255 | 22432 | 22480 | 21991 | 22047 | 22017 |
|  | cb4-3 | 22131 | Best | 22012 | 21886 | 21962 | 22131 | 21849 | 21797 | 21949 |
|  |  |  | Mean | 21662 | 21466 | 21632 | 21720 | 21313 | 21342 | 21461 |
|  | cb4-4 | 22772 | Best | 22420 | 22319 | 22463 | 22347 | 22325 | 22418 | 22376 |
|  |  |  | Mean | 22391 | 21992 | 22233 | 22160 | 21961 | 22037 | 22029 |
|  | cb4-5 | 22751 | Best | 22312 | 22440 | 22619 | 22417 | 22168 | 22215 | $22254$ |
|  |  |  | Mean | 22182 | 22132 | 22279 | 22290 | 21840 | 21822 | 21903 |
| 10x250-0.25 | cb5-1 | 59187 | Best | 58820 | 57757 | 58725 | 58476 | 56928 | 57530 | 57036 |
|  |  |  | Mean | 58812 | 56708 | 58148 | 58310 | 55759 | 55854 | 55960 |
|  | cb5-2 | 58781 | Best | 58339 | 57363 | 58321 | 57937 | 56337 | 56568 | 56490 |
|  |  |  | Mean | 58267 | 56793 | 58074 | 57790 | 55455 | 55443 | 55708 |
|  | cb5-3 | 58097 | Best | 57804 | 56690 | 57764 | 57062 | 55573 | 56426 | 55982 |
|  |  |  | Mean | 57626 | 56024 | 57372 | 56960 | 54638 | 54793 | 54727 |
|  | cb5-4 | 61000 | Best | $60597$ | 59930 | 60597 | 60326 | 58595 | 59030 | 59077 |
|  |  |  | Mean | $60460$ | 58934 | 60194 | 60030 | 57766 | 58057 | 57721 |
|  | cb5-5 | 58092 | Best | 57567 | 56863 | 57233 | 56276 | 56186 | 56217 | 56204 |
|  |  |  | Mean | 57559 | 56066 | 56961 | 56060 | 54850 | 54941 | 54872 |
| 10x500-0.25 | cb6-1 | 117821 | Best | 117371 | 113362 | 117287 | - | 108487 | 110996 | 111669 |
|  |  |  | Mean | 116882 | 112541 | 116830 | - | 105760 | 107698 | 108367 |
|  | cb6-2 | 119249 | Best | 118730 | 115022 | 118737 | - | 109569 | 114262 | 113001 |
|  |  |  | Mean | 118250 | 114250 | 118385 | - | 106775 | 108648 | 109197 |
|  | cb6-3 | 119215 | Best | 118905 | 115419 | 118488 | - | 109705 | 113987 | 112419 |
|  |  |  | Mean | 118402 | 114372 | 118003 | - | 106853 | 108576 | 109004 |
|  | cb6-4 | 118829 | Best | 117354 | 115038 | 118116 | - | 108628 | 112476 | 112198 |
|  |  |  | Mean | 116593 | 113444 | 117714 | - | 105679 | 107692 | 107796 |
|  | cb6-5 | 116530 | Best | 115334 | 112971 | 116530 | - | 106972 | 109567 | 109287 |
|  |  |  | Mean | 114447 | 111707 | 115301 | - | 104509 | 106217 | 106212 |



Figure 2 Results for the first dataset of each group of problem size

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