

On the matching polynomial of hypergraphs*

Research Article

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Abstract: The concept of the matching polynomial of a graph, introduced by Farrell in 1979, has received considerable attention and research. In this paper, we generalize this concept and introduce the matching polynomial of hypergraphs. A recurrence relation of the matching polynomial of a hypergraph is obtained. The exact matching polynomials of some special hypergraphs are given. Further, we discuss the zeros of matching polynomials of hypergraphs.

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1. Introduction

Hypergraphs are systems of finite sets. With the development of computer science, this branch of mathematics has developed rapidly during the late twentieth century. Hypergraphs model more general types of relations than graphs. For more results on hypergraph, we refer to [1, 2, 5, 11–13]. The matching polynomial of a graph was introduced by Farrell in 1979. For more details on the matching polynomial of a graph, we refer to [3, 4, 6, 7, 9, 15, 21].

An h -uniform hypergraph is a pair $H = (V; E)$, where $V = V(H)$ is a finite set of vertices, and $E = E(H) \subseteq \binom{V}{h}$ is a family of h -element subsets of V called *hyperedges*. All hypergraphs considered in this paper are h -uniform hypergraphs. A *simple hypergraph* is a hypergraph $H = (V; E)$ such that $e_i \subseteq e_j$ if and only if $i = j$. A hypergraph is *linear* if it is simple and $|e_i \cap e_j| \leq 1$, for all $e_i, e_j \in E$ where

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$i \neq j$. Given $V' \subseteq V$, the *subhypergraph* H' is the hypergraph $H' = (V', E' = (e_j)_{j \in J})$ such that for all $e_j \in E', e_j \subseteq V'$. A *matching* of H consists of isolated vertices and hyperedges only. A *k-matching* of H contains k disjoint hyperedges in $E(H)$.

Let M be a k -matching in H and let us assign “weights” w_1 and w_2 to each vertex and edge, respectively in M . Let us associate the weight $w_1^{n-kh}w_2^k$ with M . Then if a_k is the number of k -matchings in H , the total weight of the k -matchings in H will be $a_k w_1^{n-kh} w_2^k$. By summing the weights of all the k -matchings in H , for all possible values of k , we will obtain a polynomial in w_1 and w_2 . This polynomial is called the *matching polynomial* of H . Denote by it $M(H; \underline{w})$, i.e.

$$M(H; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w_1^{n-kh} w_2^k,$$

where $w = (w_1, w_2)$ is called the *weight vector* associated with the matching polynomial. If we put $w_1 = w_2 = w$, then the resulting polynomial in w will be called the *simple matching polynomial* of H ,

$$M(H; w) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w^{n-k(h-1)}.$$

The matching polynomials of hypergraphs may also be regarded as a special type of multivariate hyperedge elimination chromatic polynomials first investigated by White [20].

This paper is organized as follows. In section 2, we present a well-known fundamental theorem for the matching polynomials of hypergraphs which we use in section 3, to obtain recursively the formulae for the matching polynomials of some classes of hypergraphs. In general, these formulae are intractable, giving our choice of these particular hypergraphs which are much easier to manipulate. We hope these results can inspire further computations of these polynomials. In the last section, we discuss the zeros of these polynomials when compared to their graphs counterparts.

2. The fundamental theorem for matching polynomials

Firstly, we define two graphs operations which will be used later. Deletion of hyperedge e , i.e. e is removed and the vertices of e are still retained. Contraction of hyperedge e , i.e. e is removed and all vertices of e and hyperedges adjacent to e are removed. Since matchings consist of isolated vertices and hyperedges, the inclusion of an hyperedge in a matching implies the exclusion of all hyperedges adjacent to it. By partitioning the set of matchings in the hypergraph into two classes: (i) those containing a given hyperedge e and (ii) those not containing e , the following result is immediate.

Theorem 2.1. *Let H be a hypergraph containing a hyperedge e . Let H' be the hypergraph obtained from H by deleting e , and H'' be the hypergraph obtained from H by contracting e . Then*

$$M(H; \underline{w}) = M(H'; \underline{w}) + w_2 M(H''; \underline{w}).$$

The fundamental recurrence relation of Theorem 2.1 firstly appeared in [7], for graphs and it is widely used. Other versions of this recursion also appeared in [18].

Proposition 2.2. *Let H be a hypergraph, and let H_1 and H_2 be two subhypergraphs of H such that $V(H_1) \cap V(H_2) = \emptyset$ and $H_1 \cup H_2 = H$. Then*

$$M(H; \underline{w}) = M(H_1; \underline{w}) \cdot M(H_2; \underline{w}).$$

The *line-graph* of H is the graph $L(H) = (V'; E')$ such that:

- (i) $V' = E$ when H is without repeated hyperedge;

(ii) $\{e_i, e_j\} \in E'$ ($i \neq j$) if and only if $e_i \cap e_j \neq \emptyset$.

The first formal definition of the independence polynomial appears to be due to Gutman and Harary [10]. The independence polynomial of a graph G will be denoted by $I(G, x)$, i.e.

$$I(G; x) = \sum_{k=0}^{\alpha(G)} a_k x^k,$$

where a_k is the number of independence sets of G with exactly k vertices. For more details on the independence polynomial, we refer to [14, 16, 17, 19].

By the above definition, one can see that the coefficients of the matching polynomial of a simple hypergraph and the coefficients of the independence polynomial of its line-graph are identical. Since the *line-graph* of a hypergraph H is a simple graph, it follows that the matching polynomial of a hypergraph can be obtained by the known independence polynomial of its line-graph. The following proposition gives us a method to calculate the matching polynomial of a given hypergraph. Note that $I(L(H); x)$ is the independence polynomial of $L(H)$. Then the following proposition is immediate.

Proposition 2.3. *Let*

$$I(L(H); x) = \sum_{k=0}^{\alpha(L(H))} a_k x^k$$

be the independence polynomial of $L(H)$. Then the matching polynomial of H is

$$M(H; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w_1^{n-kh} w_2^k.$$

3. Matching polynomials of some special hypergraphs

In this section, we discuss the matching polynomials of hyperpaths, hyperstars, hypercycles, complete h -uniform hypergraphs and two new hypergraphs HW_ℓ and HS_ℓ .

3.1. Matching polynomials of hyperpaths and hyperstars

A *hyperpath* in H from x to y , is a vertex-hyperedge alternating sequence: $x = x_1, e_1, x_2, e_2, \dots, x_s, e_s, x_{s+1} = y$ such that

- $x_1, x_2, \dots, x_s, x_{s+1}$ are distinct vertices with the possibility that $x_1 = x_{s+1}$;
- e_1, e_2, \dots, e_s are distinct hyperedges;
- $x_i, x_{i+1} \in e_i$, for all $i \in \{1, 2, \dots, s\}$.

If $x = x_1 = x_{s+1} = y$, then the hyperpath is called a *hypercycle*.

Here and throughout, the length is equal to the number of hyperedges. Let P_ℓ be a linear hyperpath with length of ℓ . We can apply Theorem 2.1 to P_ℓ by operating for a terminal hyperedge. This immediately yields the recurrence relation.

Proposition 3.1. *Let P_ℓ be a linear hyperpath with length of ℓ ($\ell > 1$). Then*

$$M(P_\ell; \underline{w}) = w_1^{(h-1)} M(P_{\ell-1}; \underline{w}) + w_1^{(h-2)} w_2 M(P_{\ell-2}; \underline{w}). \tag{1}$$

Furthermore, the following expression of $M(P_\ell; \underline{w})$ can be obtained.

Theorem 3.2. *Let P_ℓ be a linear hyperpath with length of ℓ ($\ell > 1$). Then*

$$M(P_\ell; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} \binom{\ell + 1 - k}{k} w_1^{(n-kh)} w_2^k, \tag{2}$$

where n denotes the order of P_ℓ .

Proof. By Proposition 3.1, we have

$$M(P_\ell; \underline{w}) = w_1^{(h-1)} M(P_{\ell-1}; \underline{w}) + w_1^{(h-2)} w_2 M(P_{\ell-2}; \underline{w}).$$

For convenience, we use $p_{(i)}$ to denote $M(P_i; \underline{w})$. Then

$$p_{(0)} = w_1, \quad p_{(1)} = w_2 + w_1^h.$$

We can obtain the generating function for $M(P_\ell; \underline{w})$ from Equation (1). Let $g(x)$ be the generating function of $M(P_\ell; \underline{w})$.

$$\begin{aligned} g(x) &= p_{(0)} + p_{(1)}x + p_{(2)}x^2 + p_{(3)}x^3 + \dots \\ -w_1^{(h-1)} xg(x) &= -w_1^{(h-1)} p_{(0)}x - w_1^{(h-1)} p_{(1)}x^2 - w_1^{(h-1)} p_{(2)}x^3 + \dots \\ -w_1^{(h-2)} w_2 x^2 g(x) &= -w_1^{(h-2)} w_2 p_{(0)}x^2 - w_1^{(h-2)} w_2 p_{(1)}x^3 + \dots \end{aligned}$$

By Equation (1), we can obtain the generating function, i.e.

$$g(x) = \frac{w_1 + w_2 x}{1 - w_1^{(h-1)} x - w_1^{(h-2)} w_2 x^2}. \tag{3}$$

By expanding Equation (3) as the sum of an infinite geometrical progression and extracting the coefficient of x^ℓ , we can obtain the matching polynomials of hyperpath P_ℓ , i.e.

$$M(P_\ell; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} \binom{\ell + 1 - k}{k} w_1^{(n-kh)} w_2^k,$$

where n denotes the order of P_ℓ . □

Remark 3.3. *In fact, we can also prove Theorem 3.2 by the idea from Proposition 2.3. Note that the matching polynomial of a hypergraph and the independence polynomial of its line-graph are identical. Observe that the line-graph of P_ℓ is the path P with ℓ vertices. Song et al. [17] obtained the independence polynomial of the path P with n vertices, i.e.*

$$I(P_n; x) = \sum_{s=0}^{\lfloor (n+1)/2 \rfloor} \binom{n + 1 - s}{s} x^s.$$

A hyperstar S is a family of hyperedges containing a given vertex. Let S_ℓ be a linear hyperstar with ℓ edges. Since any two edges of S are intersecting, a matching of a hyperstar has at most one edge. So, we can obtain the matching polynomials of hyperstars.

Proposition 3.4. *The matching polynomial of S_ℓ is*

$$M(S_\ell; \underline{w}) = 1 + \ell w_1^{(n-h)} w_2,$$

where n denotes the order of S_ℓ .

3.2. Matching polynomials of hypercycles

Let C_ℓ be a linear hypercycle with length of ℓ . We apply Theorem 2.1 to C_ℓ by operating for an arbitrarily hyperedge. This immediately yields the recurrence relation.

Theorem 3.5. *Let C_ℓ be a linear hypercycle with length of ℓ ($\ell > 2$). Then*

$$M(C_\ell; \underline{w}) = w_1^{(h-2)} M(P_{\ell-1}; \underline{w}) + w_1^{2(h-2)} w_2 M(P_{\ell-3}; \underline{w}). \tag{4}$$

Theorem 3.6. *Let C_ℓ be a linear hypercycle with length of ℓ ($\ell > 2$). Then*

$$M(C_\ell; \underline{w}) = w_1^n + \sum_{k=1}^{\lfloor n/h \rfloor} \frac{\ell}{k} \binom{\ell-1-k}{k-1} w_1^{(n-kh)} w_2^k,$$

where n denotes the order of C_ℓ .

Proof. From Theorem 3.5, we have

$$M(C_\ell; \underline{w}) = w_1^{(h-2)} M(P_{\ell-1}; \underline{w}) + w_1^{2(h-2)} w_2 M(P_{\ell-3}; \underline{w}).$$

By Equation (2), we can obtain the expression of $M(P_{\ell-1}; \underline{w})$ and $M(P_{\ell-3}; \underline{w})$ as follows.

$$M(P_{\ell-1}; \underline{w}) = \sum_{k=0}^{\lfloor [(\ell-1) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-k}{k} w_1^{\lfloor [(\ell-1) \cdot (h-1) + 1] - kh \rfloor} w_2^k,$$

$$M(P_{\ell-3}; \underline{w}) = \sum_{k=0}^{\lfloor [(\ell-3) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-2-k}{k} w_1^{\lfloor [(\ell-3) \cdot (h-1) + 1] - kh \rfloor} w_2^k.$$

Therefore, we have

$$\begin{aligned} M(C_\ell; \underline{w}) &= w_1^{(h-2)} \sum_{k=0}^{\lfloor [(\ell-1) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-k}{k} w_1^{\lfloor [(\ell-1) \cdot (h-1) + 1] - kh \rfloor} w_2^k \\ &\quad + w_1^{2(h-2)} w_2 \sum_{k=0}^{\lfloor [(\ell-3) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-2-k}{k} w_1^{\lfloor [(\ell-3) \cdot (h-1) + 1] - kh \rfloor} w_2^k \\ &= w_1^{\lfloor [(\ell-1) \cdot (h-1) + 1] + h - 2 \rfloor} + \sum_{k=1}^{\lfloor [(\ell-1) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-k}{k} w_1^{\lfloor [(\ell-1) \cdot (h-1) + 1] - kh + h - 2 \rfloor} w_2^k \\ &\quad + \sum_{k'=1}^{\lfloor [(\ell-3) \cdot (h-1) + 1] / h \rfloor} \binom{\ell-1-k'}{k'-1} w_1^{\lfloor [(\ell-3) \cdot (h-1) + 1] - k'h + 3h - 4 \rfloor} w_2^{k'} \\ &= w_1^{\ell(h-1)} + \sum_{k=1}^{\lfloor \ell(h-1) / h \rfloor} \left(\binom{\ell-k}{k} + \binom{\ell-1-k}{k-1} \right) w_1^{\ell(h-1) - kh} w_2^k \\ &= w_1^n + \sum_{k=1}^{\lfloor n/h \rfloor} \frac{\ell}{k} \binom{\ell-k-1}{k-1} w_1^{n-kh} w_2^k \end{aligned}$$

□

Furthermore, we obtain the recurrence relation involving the matching polynomials of hypercycles only.

Theorem 3.7. Let C_ℓ be a linear hypercycle with length of ℓ ($\ell \geq 5$). Then

$$M(C_\ell; \underline{w}) = w_1^{(h-1)} M(C_{\ell-1}; \underline{w}) + w_1^{(h-2)} w_2 M(C_{\ell-2}; \underline{w}).$$

Proof. From Theorem 3.5, we have

$$M(C_\ell; \underline{w}) = w_1^{(h-2)} M(P_{\ell-1}; \underline{w}) + w_1^{2(h-2)} w_2 M(P_{\ell-3}; \underline{w}) \quad (\ell \geq 3). \tag{5}$$

By Proposition 3.1, we can obtain the following equalities:

$$M(P_{\ell-1}; \underline{w}) = w_1^{(h-1)} M(P_{\ell-2}; \underline{w}) + w_1^{(h-2)} w_2 M(P_{\ell-3}; \underline{w}) \tag{6}$$

and

$$M(P_{\ell-3}; \underline{w}) = w_1^{(h-1)} M(P_{\ell-4}; \underline{w}) + w_1^{(h-2)} w_2 M(P_{\ell-5}; \underline{w}). \tag{7}$$

By Equations (5), (6) and (7), we have

$$\begin{aligned} M(C_\ell; \underline{w}) &= w_1^{(h-1)} [w_1^{(h-2)} M(P_{\ell-2}; \underline{w}) + w_1^{2(h-2)} w_2 M(P_{\ell-4}; \underline{w})] \\ &\quad + w_1^{(h-2)} w_2 [w_1^{(h-2)} M(P_{\ell-3}; \underline{w}) + w_1^{2(h-2)} w_2 M(P_{\ell-5}; \underline{w})]. \end{aligned}$$

From Theorem 3.5, we have

$$M(C_\ell; \underline{w}) = w_1^{(h-1)} M(C_{\ell-1}; \underline{w}) + w_1^{(h-2)} w_2 M(C_{\ell-2}; \underline{w}).$$

□

Remark 3.8. In fact, we can also prove Theorem 3.6 by the idea from Proposition 2.3. Note that the coefficients of the matching polynomial of a simple hypercycle and the coefficients of the independence polynomial of its line-graph are identical. Observe that the line-graph of a cycle C_ℓ is also a cycle with ℓ vertices. Song et al. [17] obtained the independence polynomials of the cycle C with n vertices, i.e.

$$I(C_n; x) = 1 + \sum_{s=1}^{\lfloor n/2 \rfloor} \frac{n}{s} \binom{n-1-s}{s-1} x^s.$$

3.3. Matching polynomial of a complete h -uniform hypergraphs

A complete h -uniform hypergraph is a hypergraph which has all h -subsets of V as hyperedges. Let K_n^h be a complete h -uniform hypergraph. We can obtain the matching polynomial of a complete h -uniform hypergraph by applying the Principle of Combinatorial Enumeration.

Proposition 3.9. Let K_n^h be a complete h -uniform hypergraph with n vertices. Then

$$M(K_n^h; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w_1^{(n-kh)} w_2^k,$$

where

$$a_0 = 1, a_k = \frac{\prod_{i=0}^{k-1} \binom{n-ih}{h}}{k!} \quad (k \geq 1).$$

Proof. We choose h vertices from the vertex set of K_n^h to form the first hyperedge e_1 of the k -matching. Note that we have $\binom{n}{h}$ ways to do it. There are $n - h$ remaining vertices in $K_n^h - \{e_1\}$. We continue to choose h vertices from the vertex set of $K_n^h - \{e_1\}$ to form a new hyperedge of the k -matching, say e_2 . Note that we have $\binom{n-h}{h}$ ways to do it. There are $n - 2h$ remaining vertices in $K_n^h - \{e_1, e_2\}$. Continue the process by the same method. Then we can find e_1, e_2, \dots, e_k such that $e_i \cap e_j = \emptyset$ ($1 \leq i \neq j \leq k$). Since the hyperedges of a k -matching are unordered, the coefficient a_k of $M(K_n^h; \underline{w})$ is

$$\frac{\prod_{i=0}^{k-1} \binom{n-ih}{h}}{k!}$$

for the special case, $a_0 = 1$. □

3.4. Matching polynomials of HW_ℓ and its dual

A *centipede* is a tree denoted by $W_n = (A \cup B, E)$ ($n \geq 1$), where $A \cup B$ is its vertex set, $A = \{a_1, \dots, a_n\}$, $B = \{b_1, \dots, b_n\}$, and the edge set $E = \{a_i b_i : 1 \leq i \leq n\} \cup \{b_i b_{i+1} : 1 \leq i \leq n - 1\}$ (see Figure 1(a)).

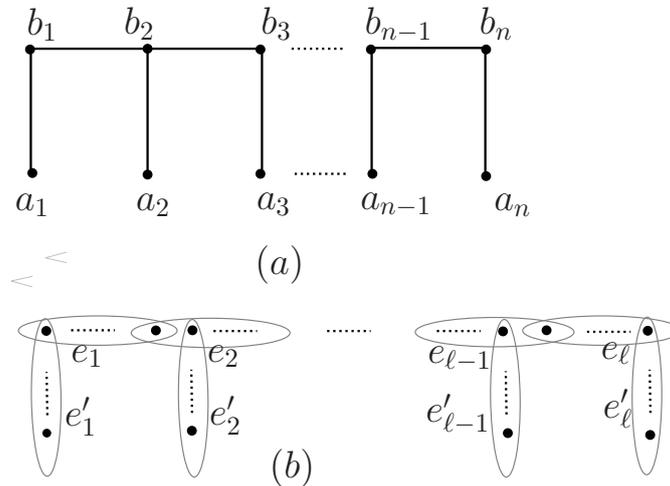


Figure 1. (a) Centipede W_n , (b) Hypergraph HW_ℓ .

We define a new hypergraph $HW_\ell = (V; E_1 \cup E_2)$ such that $e_i \in E_1$, $e'_i \in E_2$ and $e_i \cap e'_i \neq \emptyset$, where $1 \leq i \leq \ell - 1$ (see Figure 1(b)).

Lemma 3.10. [14] Let $I(W_\ell; x)$ be the independence polynomial of W_ℓ . Then

$$I(W_\ell; x) = \sum_{s=0}^{\ell} t_s x^s,$$

where

$$t_s = \sum_{j=0}^s \binom{\ell - j}{\ell - s} \binom{\ell + 1 - j}{j}, s \in \{0, 1, \dots, \ell\}.$$

Proposition 3.11. Let $M(HW_\ell; \underline{w})$ be the matching polynomial of HW_ℓ with n vertices. Then

$$M(HW_\ell; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w_1^{(n-kh)} w_2^k,$$

where

$$a_k = \sum_{j=0}^k \binom{\ell-j}{\ell-k} \binom{\ell+1-j}{j}.$$

Proof. Observe that the line-graph of HW_ℓ is a centipede W_ℓ . From Lemma 3.10, we have

$$I(W_\ell; x) = \sum_{s=0}^{\ell} t_s x^s,$$

where

$$t_s = \sum_{j=0}^s \binom{\ell-j}{\ell-s} \binom{\ell+1-j}{j}, s \in \{0, 1, \dots, \ell\}.$$

By Proposition 2.3 and the above formula, we can obtain the matching polynomial of HW_ℓ , i.e.

$$M(HW_\ell; \underline{w}) = \sum_{k=0}^{\lfloor n/h \rfloor} a_k w_1^{(n-kh)} w_2^k,$$

where

$$a_k = \sum_{j=0}^k \binom{\ell-j}{\ell-k} \binom{\ell+1-j}{j}.$$

□

Remark 3.12. According to our discussion, we observe that the coefficients of $M(P_\ell; \underline{w})$ are equivalent to those of a simple path of length ℓ . Likewise, the coefficients of $M(C_\ell; \underline{w})$ are the same as those of a simple cycle of length ℓ . However, we are unable to find an equivalent matching polynomial for hypergraph HW_ℓ , giving the importance of this results and the next one.

Let $K_\ell = (V, E)$ denote a simple complete graph with ℓ vertices. Set $V(K_\ell) = \{v_i : 1 \leq i \leq \ell\}$. Let K_ℓ^* be a simple graph obtained from K_ℓ by adding ℓ new vertices $\{u_1, u_2, \dots, u_\ell$ and the edges $\{u_i v_i : 1 \leq i \leq \ell\}$.

We now define a new hypergraph $HS_\ell = (V; E_1 \cup E_2)$ satisfying the following conditions (see Figure 2).

- (1) $E_1 = \{e_1, e_2, \dots, e_\ell\}$ and $E_2 = \{e'_1, e'_2, \dots, e'_\ell\}$;
- (2) The hypergraph induced by the edges in E_1 is a linear hyperstar;
- (3) For $e_i \in E_1$ and $e'_j \in E_2$, $e_i \cap e'_j \neq \emptyset$ for $i = j$, $e_i \cap e'_j = \emptyset$, for $i \neq j$.

Lemma 3.13. [14] Let $I(K_\ell^*; x)$ be the independence polynomial of K_ℓ^* , then

$$I(K_\ell^*; x) = (1+x)^{\ell-1} \cdot [1 + (\ell+1)x].$$

Proposition 3.14. Let $M(HS_\ell; \underline{w})$ be the matching polynomial of HS_ℓ , then

$$M(HS_\ell; \underline{w}) = (\ell-1)w_1^\ell + \sum_{k=1}^{\lfloor n/h \rfloor} \left[\binom{\ell-1}{s} + \binom{\ell-1}{s-1} \right] w_1^{(n-kh)} w_2^k.$$

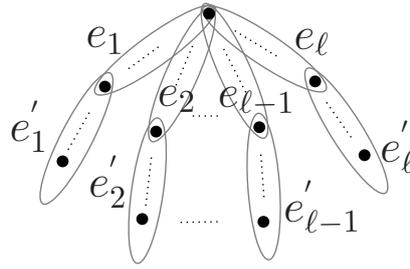


Figure 2. Hypergraph HS_ℓ .

Proof. We know that the line-graph of HS_ℓ is a K_ℓ^* . From Lemma 3.13, we have $I(K_\ell^*; x) = (1 + x)^{\ell-1} \cdot [1 + (\ell + 1)x]$. By simplifying the above equation, we obtain the following equation.

$$I(K_\ell^*; x) = (\ell - 1) + \sum_{s=1}^{\ell} \left[\binom{\ell - 1}{s} + \binom{\ell - 1}{s - 1} \right] x^s.$$

By the Propostion 2.3 and the above equation, we can obtain the matching polynomial of HS_ℓ , i.e.

$$M(HS_\ell; \underline{w}) = (\ell - 1)w_1^n + \sum_{k=1}^{\lfloor n/h \rfloor} \left[\binom{\ell - 1}{s} + \binom{\ell - 1}{s - 1} \right] w_1^{(n-kh)} w_2^k.$$

□

4. Zeros of matching polynomials of a hypergraphs

It is well known that the zeros of matching polynomial of any connected graph are all real numbers [8]. However, we observe from an example that this assertion does not hold true for hypergraphs, in general. To the best of our knowledge, this case is unknown. Consider the hypergraph H as shown in Figure 3(a). Observe that the line-graph of H is the next simple graph $L(H)$ (see Figure 3(b)).

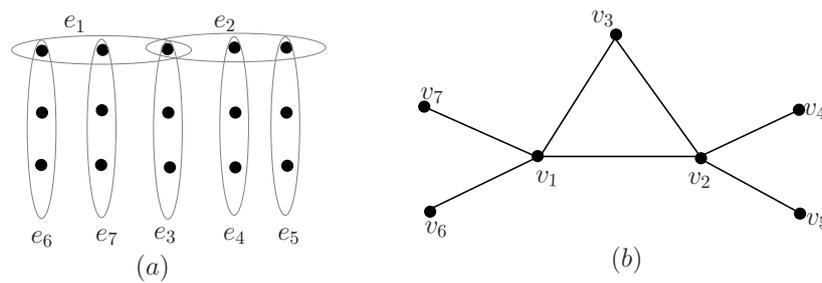


Figure 3. (a) Hypergraph H , (b) the line-graph of H .

Proposition 4.1. [14] Let $G = (V, E)$ be a graph, $U \subset V$ be such that $G[U]$ is a complete subgraph of G and $I(G, x)$ be the independence polynomial of G . Then

$$I(G, x) = I(G - U, x) + x \sum_{v \in U} I(G - N[v], x),$$

where $N(v) = \{w : vw \in E(G)\}$, $N[v] = N(v) \cup \{v\}$.

Proposition 4.2. [19] Let E_n be an edgeless graph on n vertices, $I(E_n, x)$ be the independence polynomial of E_n . Then

$$I(E_n, x) = (1 + x)^n.$$

From Proposition 4.1 and Proposition 4.2, we obtain the independence polynomial of $L(H)$, i.e.

$$I(L(H), x) = x^5 + 5x^4 + 12x^3 + 14x^2 + 7x + 1.$$

By Proposition 2.3, we obtain the matching polynomial of H , i.e.

$$M(H; \underline{w}) = w^{15} + 7w^{13} + 14w^{11} + 12w^9 + 5w^7 + w^5.$$

Using the MATLAB software, we obtain from $M(H; \underline{w})$ the following zeros:

0, 0, 0, 0, 0, -0.0000 + 2.0893i, -0.0000 - 2.0893i, -0.2839 + 0.6309i, -0.2839 - 0.6309i, 0.2839 + 0.6309i, 0.2839 - 0.6309i, 0.0000 + 1.0000i, 0.0000 - 1.0000i, -0.0000 + 1.0000i, -0.0000 - 1.0000i.

Thus, we can obtain the fact that the zeros of a matching polynomial of a hypergraph are not necessarily all real numbers.

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