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Day-ahead Management of Energy Sources and Storage in Hybrid Microgrid to reduce Uncertainty

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Highlights

• An Energy Management Algorithm is proposed for day-ahead scheduling in a microgrid.

• Optimization problem is framed to minimize the cost of power generation under DMO model.

• Artificial Fish Swarm technique is used as a tool to solve the optimization part.

· Several scenarios are drawn to treat uncertainty in Renewable Energy availability.

ArticleInfo	Abstract
Received:17/01/2019 Accepted:25/06/2019	A day ahead management strategy is proposed in this article to schedule energy generators and storage in presence of Renewable Energy Sources under uncertainty conditions with an objective to optimize the cost of energy generation. Artificial Fish Swarm algorithm is used as optimization tool. The optimization problem is framed considering all the practical constraints of energy
Keywords	generators and storage units. The uncertainty of Renewable Energy Sources is treated with a proven uncertainty model and several scenarios are drawn for energy availability and demand.
Microgrid Uncertainty Storage Energy Management Day ahead scheduling	The proposed energy management algorithm is tested numerically on a grid connected microgrid hosting a group of hybrid energy sources and storage battery for day ahead scheduling under dynamic pricing and demand side management in one of the generated uncertainty scenarios. The obtained results show that the performance of Artificial Fish Swarm algorithm as an optimizing tool is validated and the proposed Energy Management System is found to optimize the cost of energy generation while matching the power generated with power required.

1. INTRODUCTION

The economic research forecasts that the cost of energy generation from RESs will nose dive next decade in comparison with their conventional counterparts like coal [1]. In realization of these forecasts, solar generated electrical energy has witnessed a lowest bidding rate of 2.44 Rupees in India from a private producer [2]. It is comparatively much cheaper than conventional sources like nuclear or coal. Apart from cost of energy generation, the potency of solar generated energy in meeting the global energy needs with almost zero damage to the environment is yet another advantage. The sunshine rich parts of the world like Asia, Africa, and Australia etc have enough solar energy sources to meet a major share of global energy requirements. The idea of tapping solar and wind energy which are distributed sources by availability has encouraged the concept of connecting smaller and distributed generation sources to the utility grid at tail ends (distribution systems) at low and medium voltages levels. This lead to a whole new concept of autonomous microgrids [3]. The microgrids operate in either grid connected mode or in autonomous mode. The developments in microgrid enabling technologies like capacity building of DG units from a few kW to MW, efficient energy derivation methods from RESs, reducing costs of energy storage, availability of wide area controlling technologies like SCADA along with incentives offered by governments for renewable energy generation have grown so attractive that they can no more be overlooked by investment policy makers. The investment policies in energy market are being reframed to accommodate microgrids into the existing utility grid system and symptoms that conventional centralized generation concept of energy is slowly shifting towards the distributed generation concept are evident [4].

All the above advantages being on one side of the microgrids, they suffer from quite a few operation and control difficulties. Uncertainty of energy availability from the RESs is one such serious setback, which ultimately affects the reliability of the system, if not treated properly. The energy output from a RES generator is not under the control of the microgrid operator. He can only curtail the generation in the event of excess generation. A lot of research is inspired by the need to reduce the uncertainty of energy availability from RESs and to balance the generation with the loads in the microgrids. The classical energy balance techniques fall short in microgrid environment due to the uncertainty in generation and loads added with bi-directional power flows, which add even more difficulty in energy balancing. Thanks to research, it has come out with quite a few modern techniques modelled and tested to contain the uncertainty. Day ahead scheduling using data mining, maintaining battery storage back-ups and spinning reserves, demand side management tools like load shifting, load curtailing (load shedding under deficit generation), demand response etc top the list. A sample of literature using the above energy balance techniques to solve various types of objective functions include mixed integer linear programming problems [5-7], linear programming problems [8]. Different types of linear programming techniques and nature inspired swarm intelligence algorithms are employed to solve the problems of optimum energy exchange (charging/discharging) schedules of batteries/storage facility [9-11]. All the above energy balancing techniques are from the generation side. The other energy balancing techniques from demand side are also tested successfully, like load shifting and shedding [12], demand response [13-16]. Load shedding being a technique used to manage the constant power loads (They are either supplied with rated power or else completely disconnected), using electric springs is a new technique to supply constant energy loads with reduced power [17]. Apart from this quite a large number of models are proposed and tested to represent the uncertainty of energy availability with the RES generators and the loads [18-21].

The problem of scheduling power generation sources and storage devices while optimizing the cost of power generation from sources is considered in this article. The solution to the problem needs two aspects to be solved. One is to find the optimal mix of generation from different sources to minimize the cost of generation and the other part is to schedule the required power among generation sources following the optimized mix as evaluated. Until late, the cost of power generation from renewable sources is considered as zero on the pretext that they need zero expenditure towards the fuel and hence they are treated as nondispatchable sources [22,23]. This assumption is too remote from practical reality and simplifies the optimization problem too much. A few works reported in literature have considered the cost of power generation from renewable sources as non-zero. Authors in [24] have considered cost of generation from wind and PV units as quadratic functions. Authors in [25] have proposed a new cost function for renewable power generation by computing a penalty amount for differential power between actual powers generated and forecasted power. This case is specifically proposed for the microgrid operators who do not have sufficient facilities to store all the excess energy available with renewable sources. The penalty amount so computed is made a component in computing the cost of renewable energy generation. The authors in [26] put an end to general flow of concept in literature that the power generation from renewable sources in uncontrollable. They have proposed an Automatic Generation Control strategy to control power generation from wind farms.We have considered both the conventional and renewable generators in dispatchable category with respective cost functions for both. The popular optimization techniques like GA and PSO, which are the favourites in literature, suffer from serious drawbacks in terms of computational complexities and time. More constrained the problem is, more will be these complexities. For instance, the process of fine tuning the computational parameters like fitness normalization, mutation rate, crossover parameters etc in GA is more trial and error based. Such complexities are quite lesser in number in Artificial Fish Swarm (AFS) algorithm. Though it suffers from the disadvantage of not using the past experience of the swarm members in evaluating the next step in search, AFS algorithm has the outweighing merits of higher accuracy, lesser computational time, faster convergence, very few parameters to be tuned etc. As such, AFS algorithm is used in this paper to find solution for the optimization part of the problem. The problem is framed with an objective to reduce the cost of energy generation in presence of uncertain energy availability from the RES generators under dynamic pricing of energy. A bunch of renewable and non-renewable energy sources are considered along with battery back-up. The uncertainty with RESs energy availability and loads is treated using a proven uncertainty model. The battery is charged from the excess energy available and is used to back-up supply under deficit generation. The microgrid is operated in grid connected mode with an upper limit for power drawl. The problem formulation is done with a view to

include demand side management which allows load shedding under power shortage condition, compensating the consumer for the inconvenience caused due to power disconnection. Artificial Fish Swarm Algorithm is employed to work out the optimization part of the objective function. The problem is solved over hourly intervals for 24 hours.

2. ARTIFICIAL FISH SWARM ALGORITHM

The fishes in a swarm are seen continuously changing their positions. For an observer, their movements appear random and unsynchronized. But in reality, the fishes do it purposefully and with a high sense of synchronization. They move continuously to adjust their position in synchronization to the path leading to a location of their objective. Their objective could be anything like staying protected from a predator's attack, avoiding collision with neighbouring fellow fishes, finding a location where chances of finding food is better etc. Such movements of the fishes are modelled and the AFS algorithm is framed [27]. It frames a localized searching method for each of the individual artificial fish and each search ends with updating the solution to obtain the best location for the objective chosen. Each individual fish in the swarm is termed as artificial fish. Each instantaneous position that these artificial fishes occupy on their path to the objective location, symbolizes a possible solution to the problem at hand. These instantaneous positions are modelled differently for different objectives of movement by fish. The fishes exhibit different movement patterns for different behaviours like random searching behaviour, trail following behaviour, preying, swarming and mating behaviours etc. The AFS algorithm adopts these entire behaviours one after the other in its quest for the optimum solution. The AFS algorithm can be summed up as follows.

Let the current location of an individual artificial fish be

$$X_i = (x_{i1}, x_{i2}, \dots, x_{in}), \tag{1}$$

where i signifies the number of control variables and n signifies the total number of individual artificial fishes living in the swarm.

Then the consistency of fulfilling its objective at the location will be

$$Y_i = f(X_i). (2)$$

On its way to the location where the consistency of fulfilling objective is best, the artificial fish alters its position to a new location X_i^* which is given by

$$X_{i}^{*} = X_{i} + rand() * step * \frac{X_{j} - X_{i}}{\|X_{j} - X_{i}\|},$$
(3)

where *rand* () is any random number that lies between '0' and '1', *step* is the largest radial displacement through which the fish can navigate in one move and X_j is anyone of the calculated positions which exists within the visual extent of the fish. The calculated position X_j is a function of behaviour of the artificial fish and it is different for different behaviours.

- 1. Chasing behaviour: In this behaviour, the artificial fish simply follows the fish in its front. Condition is that the position of the neighbouring fish exists within its visual extent. In this behaviour, the position X_j is defined as
 - $X_j = The position of the neighbour fish.$
- Swarming / gathering behaviour: In this behaviour the artificial fish tries to move close to the centre of the swarm to ensure that it is surrounded by the fellow fishes in all directions to avoid any potential attack/predator from outside. In this behaviour, the position X_j is defined as X_j = X_i + rand () * step * ^{X_c-X_i}/<sub>||X_c-X_i||, (4)
 </sub>

where the geometrical centre location of the swarm is denoted by X_c .

3. Foraging / Preying behaviour: In this behaviour, the artificial fish estimates the probability of finding food at different locations either by vision or sense. It searches for a location with better food consistency and directly moves in that direction. The position X_j in this behaviour is defined as

$$X_i = X_i + rand() * visual.$$



Figure 1. Flow chart of AFS algorithm

The AFS algorithm uses all the above explained behaviours of fish movement in a sequential order in search for the optimality of the objective function. With each step, the better location is updated and the search is continued until the objective is achieved. Figure-1 shows the flow chart of the AFS algorithm. After initialization of the parameters, the algorithm generates the initial fish swarm for generator x_i for i=1 to n as shown in Figure-1. The different artificial fishes are tested for the fitness of the objective function. The various behaviours of the fish swarm like chasing the trail, swarming and foraging behaviours, are simulated in a sequence on the artificial fishes generated as shown in Figure 1 and the best resulting locations are updated with each iteration. If the termination condition is true, the simulation is stopped.

The parametric mapping between the AFS algorithm and the problem considered is done as follows. The power generation to be scheduled to each generator is mapped with artificial fish, the vision of the fish is mapped with maximum power limit of the generator and the step of the fish is mapped with generation ramp allowed with the individual generators.

3. MODELLING OF RELEVANT AGENTS

The mathematical models of different agents used in this problem such as generators, storage devices, cost of energy, uncertainty associated with the loads and renewable energies availability etc are detailed as below.

3.1. Generator model

Let $P_i^{con}(t)$ and $P_i^{res}(t)$ be the powers generated by i^{th} conventional generator (non-renewable energy)

and j^{th} renewable generator at time interval t. The upper and lower limits of their power generation are given by

$$P_i^{con_{min}} \le P_i^{con}(t) \le P_i^{con_{max}}, t \in T, i \in I,$$
(6)

$$P_i^{res_{min}} \le P_j^{res}(t) \le P_j^{res_{max}}, t \in T, j \in J.$$
(7)

3.2. Storage model

Battery storage is modelled as a distributed unit and their charging and discharging schedule is controlled by the energy management system based on the following constraints.

Let $P_k^{stg_{ch}}(t), t \in T, k \in K$ and $P_k^{stg_{dch}}(t), t \in T, k \in K$, be the charging and discharging powers respectively.

Then the rates of charging and discharging are bound by upper and lower limits as in (8) and (9)

$$0 \le P_k^{stg_{ch}}(t) \le P_k^{stg_{ch}}, t \in T, k \in K,$$
(8)

$$0 \le P_k^{stg_{dch}}(t) \le P_k^{stg_{dch}}, t \in T, k \in K,$$
(9)

$$SoC_k^{min} \le SoC_k(t) \le SoC_k^{max}, t \in T, k \in K,$$
(10)

$$SoC_k(t+1) = SoC_k(t) + \eta_k P_k^{stg_{ch}}(t), t \in T, k \in K,$$
(11)

$$SoC_k(t+1) = SoC_k(t) - \eta_k P_k^{stg_{dch}}(t), t \in T, k \in K.$$

$$(12)$$

The magnitude of energy stored in the battery storage facility is expressed in terms of state of charge, $SoC_k(t)$ which could either be in terms of percentage of full charge or in terms of kWh remaining for backup. It is convenient to use kWh remaining for back-up in the present case. The state of charge is bound by upper and lower limits (10). The state of charge at any instant is given by (11) while charging and by (12) while discharging, inclusive of charging and discharging efficiencies as noted in equations.

3.3. Load model

A constant power model is used for loads. In such models the loads are an aggregate of rated powers, which should be supplied in full. A controllable/load shedding version of the model allows the energy management system to either supply the load fully or disconnect it from the grid. The load in excess to total power available with the microgrid cannot be supplied and is disconnected/ shed. Such load is termed as load shed $L_m^{shd}(t)$ [28]. The load shedding is allowed with compensation $C_m^{shd}(t)$ due to the energy user by the utility operator for the inconvenience caused.

The Distribution Market Operator (DMO) model [29], which is used in the present work, restricts power transactions between the utility grid and microgrid to a predetermined award, which is communicated to the microgrid operator by the DMO one day in advance. The aim of this restriction is to reduce penetration of uncertainty from microgrids into the utility grid. This award is binding on the microgrid operator and he has to draw power from the utility grid without deviating from this award. Any deviation from the award will be penalised. In other words, the DMO decides the upper and lower limits for power drawl from the utility grid is no more an infinite source. Under such conditions, if the load to be supplied by the microgrid is more than the sum of all internal sources, storage and maximum grid power limit, the excess load is disconnected and is termed as load shed.

3.4. Modelling of costs of energy

The business operation of microgrid involves various components towards the calculation of cost of energy production. Such components which are considered in the present article may be summed up as follows.

$$C_{i}^{con}(t) = X_{i}^{con}(t) * c_{i}^{con}(t) * P_{i}^{con}(t) + U_{i}^{con}(t) * c_{i}^{con_{start}} + V_{i}^{con}(t) * c_{i}^{con_{shut}}, i \in I, t \in T,$$
(13)

$$C_{j}^{res}(t) = X_{j}^{res}(t) * c_{j}^{res}(t) * P_{j}^{res}(t) + U_{j}^{res}(t) * c_{j}^{res_{start}} + V_{j}^{res}(t) * c_{j}^{res_{shut}}, j \in J, t \in T,$$
(14)

$$C_k^{stg}(t) = Y_k^{stg}(t) * c_k^{stg} * \eta_k * P_k^{stg}(t),$$
(15)

$$C_m^{shd}(t) = Y_m^{shd}(t) * c_m^{shd}(t) * L_m^{shd}(t),$$
(16)

$$C^{grd}(t) = c^{grd}(t) * P^{grd}(t).$$
⁽¹⁷⁾

The hourly cost of energy generation $C_i^{con}(t)$ by the conventional energy source including the startup and shutdown costs is given by (13) and that of RES generator $C_j^{res}(t)$ is given by (14) which include their startup and shutdown costs. U and V are state vectors which give the shut-down or start-up status of the sources. A '1' is used to indicate if the generator is either started or shut down at the time interval and a '0' is used to indicate if the generator is already OFF state or is already in ON state.

The cost of power exchange from battery/storage device $C_k^{stg}(t)$ is modelled as given by (15), where $Y_k^{stg}(t)$ is a state vector signifying the type of power exchange of the battery/storage. $Y_k^{stg}(t)$ models the three types power exchange of the k^{th} battery/storage with the microgrid at time t. The three states are charging under excess generation, discharging under deficit generation and idle states. The energy management system schedules battery charging under excess generation and discharging under deficit generation. The term η_k refers to efficiency while charging and discharging of the battery/storage facility.

The cost of load shedding by the microgrid operator is given by (16). $Y_m^{shd}(t)$ is a state vector signifying ON/OFF state of the m^{th} load at time t. A '0' in the vector signifies an ON state and a '1' signifies OFF state. Compensation cost adds to the total cost of microgrid operation in case the load shedding is done by the microgrid operator. The term $c_m^{shd}(t)$ is the rate of compensation payable and $L_m^{shd}(t)$ is the load magnitude disconnected from the microgrid by the utility operator.

The microgrid is operated in grid connected mode. Under DMO model, the power drawl from the grid is subject to upper and lower limits and the cost of power drawl $C^{grd}(t)$ is laid down by (17). The grid power comes at a dynamic price of $c^{grd}(t)$ per unit energy.

3.5. Modelling uncertainty of RES generation

Let $P_j^{res}(t)$ be the power developed by j^{th} RES generator at time interval t and vector Ψ_j be a vector that collects all the values of $P_j^{res}(t)$ for $t \in T$. Then (18) consolidates this relation

$$\Psi_j = \left[P_j^{res}(1), P_j^{res}(2) \dots P_j^{res}(T) \right], \text{ over the time horizon } t \in T.$$
(18)

To model the uncertainty of power $P_j^{res}(t)$ over the time period $t \in T$, it is assumed that the value of Ψ_j is unknown and lies in an uncertainty set Ψ . Let there be an exclusive uncertainty set for each of renewable generators over the time period $t \in T$ and another exclusive uncertainty set for the total renewable power $\sum_{j,s} P_j^{res}(s)$ over a time sub-horizon $s \in T$ and say all these uncertainty sets form the subsets of the main uncertainty set Ψ . Then (19) relates all the above uncertainty sets,

$$\Psi = \begin{cases} \left(\left[P_j^{res}(t) \right] \middle| P_j^{res_{min}}(t) \le P_j^{res}(t) \le P_j^{res_{max}}(t), |t \in T, j \in J \right), \\ \left(\left[\sum_{j,s} P_j^{res}(s) \right] \middle| P_j^{restotal_{min}}(s) \le \sum_{j,s} P_j^{res}(s) |s \in T \le P_j^{restotal_{max}}(s) \right) \end{cases},$$
(19)

 $P_j^{res_{min}}(t)$ and $P_j^{res_{max}}(t)$ are the minimum and maximum bounds for the values of power available with the *j*th RES sourced at time interval *t*, $P_j^{restotal_{min}}(s)$ and $P_j^{restotal_{max}}(s)$ represent the minimum and maximum bounds for the values of the total energy available with the RES generators over a time subhorizon *S*. The numeric values of $P_i^{res}(t)$ and $\sum_{j,s} P_j^{res}(s)$ lie always within these respective bounds.

This uncertainty model is quite general in nature and it takes into consideration the climatic parameters of the RES sources as the model is based on the recorded data. This is a simple model in terms of computational complexity and is quite proven practically as it is based on the actual historical data. It

requires only the values of $P_j^{res_{min}}(t)$, $P_j^{res_{max}}(t)$, $P_j^{restotal_{min}}(s)$ and $P_j^{restotal_{max}}(s)$. Numeric values of these terms can be estimated from historical data analysis using any of the data inference techniques reported in literature [3].

4. MICROGRID PROBLEM FORMULATION

Minimization of energy generation cost being the prime objective of the microgrid utility operator, it can be stated as in (20)

$$\min \sum_{n=1}^{N} \sum_{t=1}^{T} f[P_n(t)],$$
(20)

where

$$f[P_{n}(t)] = \begin{cases} \sum_{t=1}^{T} \sum_{i=1}^{I} \left[W_{i}^{con}(t) * c_{i}^{con}(t) * P_{i}^{con}(t) + x_{i} * c_{i}^{conup} + y_{i} * c_{i}^{condown} \right] \\ + \sum_{t=1}^{T} \sum_{j=1}^{J} \left[X_{j}^{res}(t) * c_{j}^{res}(t) * P_{j}^{res}(t) + x_{j} * c_{j}^{resup} + y_{i} * c_{j}^{resdown} \right] \\ + \sum_{t=1}^{T} \sum_{k=1}^{K} \left[Y_{k}^{stg}(t) * c_{k}^{stg}(t) * \eta_{k} * P_{k}^{stg}(t) \right] \\ + c^{grd}(t) * P^{grd}(t) + \sum_{t=1}^{T} \sum_{m=1}^{M} \left[Z_{m}^{shd}(t) * c_{m}^{shd}(t) * L_{m}^{shd}(t) \right], \\ i, j, k, m \in n, i \in I, j \in J, k \in K, m \in M, t \in T \end{cases}$$

$$(21)$$

subject to the following constraints.

4.1. Power balance constraints

 $P^{generated}(t) = P^{required}(t),$

where

$$P^{generated}(t) = \sum_{i=1}^{I} P_i^{con}(t) + \sum_{j=1}^{J} P_j^{res}(t) + \sum_{k=1}^{K} P_k^{stg_{dch}}(t) + P^{grd}(t),$$

$$P^{required}(t) = \sum_{m=1}^{M} [L_m(t) - L_m^{shed}(t)] + \sum_{k=1}^{K} P_k^{stg_{ch}}(t), i \in I, j \in J, k \in K, m \in M, t \in T.$$
(22)

Equation (22) is the power balance equation.

4.2. Generator constraints

$$P_{i,t_{min}}^{con} \le P_{i,t}^{con} \le P_{i,t_{max}}^{con}, i \in I, t \in T,$$
(23)

$$P_{j,t_{min}}^{res} \le P_{j,t_{max}}^{res}, j \in J, t \in T.$$

$$(24)$$

Equation (23) refers to minimum/maximum power generating limits of conventional generators whereas (24) refers to that of RES generators.

4.3. Storage facility constraints

The limits of battery/storage power exchange with the microgrid while charging and discharging are given by (25)

$$0 \le P_k^{stg_{ch}}(t) \le P_k^{stg_{ch_{max}}}, 0 \le P_k^{stg_{dch}}(t) \le P_k^{stg_{dch_{max}}}, k \in K, t \in T.$$
(25)

The battery is charged from excess energy available. This constraint for the charging of the battery may be modelled as (26)

$$P_{k}^{stg_{ch}}(t) > 0 | \{ P_{i}^{con}(t) + P_{j}^{res}(t) + P^{grd}(t) \} > [L_{m}(t) - L_{m}^{shed}(t)] \text{ and } P_{k}^{stg_{ch}}(t) = 0 | else$$
(26)

$$SoC_{k,t}^{\min} \le SoC_k(t) \le SoC_{k,t}^{\max}, t \in T, k \in K.$$
(27)

Equation (27) restricts the state of charge of the battery/storage facility.

4.4. Grid power drawl constraints

 $P^{grd}(t) \leq P^{grd_{max}}$.

(28)

Equation (28) imposes limits on power drawl from the grid and restricts it to the DMO imposed award, as explained in section 3.3.

5. ALGORITHM FOR PROPOSED ENERGY MANAGEMENT SYSTEM

The proposed algorithm for day ahead energy management is implemented in following steps.

- 1. Read data and initialize.
- 2. Generate scenarios of uncertainty for RES generators and loads.
- 3. For each scenario
 - a. If total generation available is more than loads
 - i. Find power required = $load + P_{charge}$.
 - ii. Implement AFS algorithm and schedule generation.
 - iii. Charge the battery with excess power available and update state of charge.
 - b. If the total generation is less than loads
 - i. Find load shedding = load (Sum of all generation available+ $P_{dischrge}$)
 - ii. Implement load shedding.
 - iii. Find power required = load-loadshedding
 - iv. Use all the available generation
 - v. Discharge the battery to meet load and update state of charge
- 4. Evaluate cost of generation
- 5. Save data and end.

The objective of the proposed Energy Management Algorithm (EMA) is to schedule the generation and storage in a microgrid to reduce the effect of uncertainty in RES power availability and loads. Optimization of cost of generation is only a part of proposed EMA. The proposed EMA calls the Artifical Fish Swarm (AFS) algorithm when there is a chance to optimize the cost of generation in step 3.a. Optimization is possible only during excess generation, i.e., when the generation is more than the load. On the other hand, during deficit generation i.e., when the generation is less than the load, all the sources should be used and hence there is no chance for optimization. After using all the generators to the full capacity, the EMA uses storage and load shedding as options to match the load with generation in step 3.b of the proposed algorithm.

6. NUMERICAL EXAMPLE

To test the performance of the proposed EMS algorithm, a microgrid with a mix of definite energy sources (one fuel cell, one micro turbine and grid connection with upper limit of power drawl- whose maximum power availability is defined) and uncertain energy sources (one wind turbine and two photovoltaic panels) is considered. The details of the mentioned power generators are as tabulated in Table 1. The forecast of the maximum power available with the wind turbine, photovoltaic panel and the loads are as plotted in Figure 2. The sampling and scheduling is done on hourly basis for 24 hours. The dynamic prices of energy generation from fuel cell, micro turbine and the utility grid are as tabulated in Table 2. A battery of 300 kW, which is optimum [30], approximately 15% of the microgrid capacity, is considered for back-up, whose technical details are as shown in Table 1. The DMO imposed upper limit for grid power drawl is fixed at 300 kW.

▲.	uble 1. Teenmean Denands of the generators used (ktt)								
	Туре	Capacity	Max power limit	Min power limit	Cost of energy (Euro/kW)				
	Wind tubine	800	800		0.0307				
	PV panels, 1	200	200		0.0307				
	Fuel cell	500	500	30	As per Table-2				

Table 1. Technical Details of the generators used (kW)

Micro turbine	500	500	30	As per Table-2
Grid		300		As per Table 2
Battery	300	100	Max SoC	
			Min SoC	



Figure 2. Forecasted power availability of RES generators and load

Hour	Pgrid	P _{mt}	P _{fc}	Hour	Pgrid	P _{mt}	P _{fc}
1	0.027	0.0823	0.1277	13	0.572	0.0885	0.1308
2	0.02	0.0831	0.1285	14	0.286	0.0885	0.138
3	0.017	0.0831	0.129	15	0.279	0.09	0.1315
4	0.017	0.0838	0.1285	16	0.086	0.0908	0.1331
5	0.029	0.0838	0.1292	17	0.059	0.0915	0.1331
6	0.033	0.0846	0.1292	18	0.05	0.0908	0.1338
7	0.054	0.0854	0.13	19	0.061	0.0885	0.1331
8	0.215	0.0862	0.1308	20	0.181	0.0862	0.1315
9	0.572	0.0862	0.1315	21	0.077	0.0846	0.1308
10	0.572	0.0892	0.1323	22	0.043	0.0838	0.13
11	0.572	0.09	0.1315	23	0.037	0.0831	0.1285
12	0.215	0.0885	0.1308	24	0.035	0.0832	0.1281

Table 2. Cost of energy in Euros/kWh

To treat the uncertainty in maximum power availability from RES generators and hourly loads, several scenarios are generated using the uncertainty model explained in (19) and presented graphically in Figure 3 to Figure 6. The proposed energy management algorithm is implemented in scenario 3 using AFS algorithm and the results are tabulated in Table 4. The numeric data of scenario 3 is presented in Table 3 for easy comprehension. To initialize the parameters, the scheduled power of each generator (artificial fishes) is set to minimum power limit according to Table 1, the vision of the fishes is set to the maximum power limit of the generator and the step is set to 1 kW in common to all the generators. The results obtained using AFS algorithm are validated by comparing with the results obtained for implementing the same problem using AIMD algorithm. The results using AIMD algorithm are shown in Table 5.



Figure 3. Scenarios for Max Power available with Figure 4. Scenarios for Max Power available with PV1 PV2



Figure 5. Scenarios for Max Power available with wind turbine

Figure 6. Scenarios for loads

18 20 22 24

Table 3. P	ower availabili	y and load in	uncertainty	v scenario 3 (kW)
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		i i entrene i				9 800.000	10 0 (· /	
Hour	Pwind	P _{pv1}	P _{pv2}	load	Hour	Pwind	P _{pv1}	P _{pv2}	Load
1	415	0	0	1352	13	198	125	160	1721
2	231	0	0	1162	14	79	125	136	1968
3	39	0	0	1069	15	19	124	140	1837
4	20	0	0	1192	16	116	70	80	1793
5	155	0	0	1149	17	324	46	39	2035
6	412	0	15	1268	18	617	0	1	2590
7	479	43	74	1497	19	756	0	0	2096
8	518	90	90	1392	20	572	0	0	2358
9	464	103	120	1644	21	870	0	0	2676
10	422	118	165	1332	22	786	0	0	1935
11	432	157	160	1702	23	825	0	0	1637
12	332	144	196	1699	24	865	0	0	1649
Total e	energy a	vailabi	lity ove	9946	1145	1376			

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Hour	Pgwind	Pg _{pv1} .	Pg _{pv2}	Pg _{fc}	Pg _{mt}	Pgrid	P_{disch}	Cost	P_{charge}	SoC	Loadshed
1	415	0	0	237	500	300	0	92.26	100	130	0
2	231	0	0	220	481	300	0	81.33	70	230	0
3	36	0	0	240	500	293	0	78.6	0	300	0
4	5	0	0	399	488	300	0	97.42	0	300	0
5	151	0	0	224	499	275	0	83.37	0	300	0
6	389	0	8	153	448	270	0	78.77	0	300	0
7	472	8	70	189	477	281	0	97.36	0	300	0
8	495	79	54	205	499	60	0	102.01	0	300	0
9	423	51	116	498	488	68	0	164.56	0	300	0
10	415	102	106	230	470	9	0	96.63	0	300	0
11	412	126	132	459	489	84	0	172.99	0	300	0
12	312	144	119	440	491	192	0	160.15	0	300	0
13	198	125	160	500	500	238	0	260.61	0	300	0
14	79	125	136	500	500	300	100	209.49	0	300	228
15	19	124	140	500	500	300	100	203.14	0	200	154
16	116	70	80	500	500	300	70	145.92	0	100	157
17	324	46	39	500	500	300	0	142.56	0	30	326
18	617	0	1	500	500	300	0	146.27	0	30	672
19	756	0	0	500	500	300	0	152.31	0	30	40
20	572	0	0	500	500	300	0	180.71	0	30	486
21	870	0	0	500	500	300	0	157.51	0	30	506
22	786	0	0	486	463	300	0	139.01	100	30	0
23	820	0	0	137	480	300	0	93.77	100	130	0
24	836	0	0	160	423	300	0	91.85	70	230	0

Table 4. Scheduling results for scenario 3 using AFS algorithm (power in kW and cost in Euros)

Table 5. Scheduling results for scenario 3 using AIMD algorithm (powers in kW and cost in Euros)

Hour	Pgwind	Pg_{pv1} .	Pg _{pv2}	Pg_{fc}	Pg _{mt}	P _{grid}	P_{disch}	Cost	P _{charge}	SoC	Loadshed
1	364	0	0	394	394	300	0	102.01	100	130	0
2	231	0	0	350.5	350.5	300	0	87.26	70	230	0
3	39	0	0	365	365	300	0	83.71	0	300	0
4	20	0	0	436	436	300	0	98.28	0	300	0
5	155	0	0	347	347	300	0	87.37	0	300	0
6	298.25	0	15	328.25	328.25	298.25	0	89.64	0	300	0
7	340	43	74	370	370	300	0	109.93	0	300	0
8	288	90	90	318	318	288	0	145.29	0	300	0
9	353.68	103	120	383.67	383.67	299.98	0	272.82	0	300	0
10	247.26	117.98	165	277.25	277.25	247.25	0	219.12	0	300	0
11	341.68	157	160	371.67	371.67	299.98	0	274.13	0	300	0
12	332	144	196	363.5	363.5	300	0	167.85	0	300	0
13	198	125	160	469	469	300	0	289.28	0	300	0
14	79	125	136	500	500	300	100	209.49	0	300	228
15	19	124	140	500	500	300	100	203.14	0	200	154
16	116	70	80	500	500	300	70	145.92	0	100	157
17	324	46	39	500	500	300	0	142.56	0	30	326
18	617	0	1	500	500	300	0	146.27	0	30	672
19	756	0	0	500	500	300	0	152.31	0	30	40
20	572	0	0	500	500	300	0	180.71	0	30	486

21	870	0	0	500	500	300	0	157.51	0	30	506
22	735	0	0	500	500	300	0	142.36	100	30	0
23	459	0	0	489	489	300	0	128.66	100	130	0
24	453	0	0	483	483	300	0	126.47	70	230	0



Figure 7. Scheduling results using AFS algorithm



Figure 8. Scheduling results using AIMD algorithm



Figure 9. Comparison of cost of generation using AFS and AIMD algorithms





Figure 11. State of charge of battery

7. RESULTS AND DISCUSSIONS

The results of generation scheduling by implementing proposed algorithm for scenario-3 is presented in Table 4 and that using AIMD algorithm in Table 5. The term Pg with suffix in the Tables 4 and 5 refers to the power generation scheduled to the respective generators after implementing the proposed Energy Management Algorithm. In either case, the fundamental constraint of energy balance as laid down by (19) is verified and the power produced (Power produced=Power generated by conventional generators+Power generated by renewable generators+Power obtained by battery discharge+Power drawn from grid, (22)] is exactly matching the power required [Power required = Total load on microgrid-Load shed+Powre required to charge battery, (22)] at each hour. A close comparison of Table 4 and Table 5 shows that AFS algorithm is far better than AIMD algorithm in optimizing the cost of generation. For example at hour 1, the cost of grid energy is the cheapest, followed by wind and PV power, micro turbine and fuel cell in sequence of increasing cost of generation. During this hour, the proposed EMS using AFS algorithm is able to schedule complete power available from cheaper sources (grid, wind turbine, PV panels, micro turbine) and use the costliest source, i.e., fuel cell very sparingly. Whereas AIMD algorithm is unable to do the same. It is scheduling equal powers to all the generators subject to min and max powers. For micro turbine and fuel cell, the min powers are 30 kW (Table 1) and for wind turbine and PV panels, the min powers are zero. Max powers available for wind turbine and PV panels are 415 kW,0 kW,0 kW (Table-3, uncertainty scenario 3) during hour 1 and for micro turbine, fuel cell and grid it is 500 kW,500 kW and 300 kW respectively (Table 1). At hour 1, AIMD algorithm is allocating equal power scheduling to wind (364 kW), fuel cell (394-30 (min power) = 364 kW) and micro turbine (394-30 (min power) = 364 kW) and 300 kWto grid as its maximum power limit is 300 kW. In other words, the AFS algorithm is able to choose the cheaper sources and schedule more load to them compared to the costlier sources, whereas AIMD algorithm is scheduling equal amounts of loads to all the sources available irrespective of their cost of energy generation and that is the reason why there is a difference in cost of energy generation estimated by the two algorithms. Similar scheduling is done for hours 1-13 and hours 22-24. During these hours, the demand is less than the total power availability and hence the AFS algorithm is able to search for the most economic mix of power generation from different sources. During the hours 14-21, the demand is more than the total power available, and hence both the algorithms are using entire power available, there is no choice of selection of generators for scheduling and hence the cost estimation by both the algorithms is same. The above discussion compares the performance of two tools used for optimization, i.e., AIMD algorithm and AFS algorithm. Figure 7 and Figure 8 show the scheduling results using two algorithms. Figure 9 compares the costs of energy generation estimated by the two optimization tools. As far as the performance of the Energy Management System proposed, it is able to balance the power required with the power generation at each hour. The battery is assumed to be at min state of charge (30 kW) at the beginning of hour 1. The EMS is able to charge the battery to its maximum state of charge without violating the maximum rate of charge, i.e., 100kW, during the hours when excess generation is available after supplying the load. Similarly it is able to discharge the battery when required, i.e., during hours 14,15 and 16 without violating the maximum discharge rate, i.e., 100 kW. Figure 10 shows the hourly power exchange of the battery on hourly basis. The proposed EMS is also able to maintain the state of charge of the battery within the upper and lower limits as mentioned in Table 1. Figure 11 shows the hourly state of charge of the battery over the day. The EMS is also able to implement load shedding during deficit generation hours, after discharging the battery. During the peak hours, i.e., between hours 14-21, the EMS is implementing load shedding to balance the power required with power generated.

8. CONCLUSIONS

An energy management system for day ahead energy planning in a microgrid is proposed in this article with an objective to optimize the energy generation cost and to reduce the effect of uncertainty due to renewable sources. The proposed EMS is tested on a numerical problem in a grid connected microgrid with several hybrid sources like a fuel cell, a micro turbine, PV and wind generators along with battery backup under dynamic pricing policy. Several uncertainty scenarios are drawn for power availability from renewable generators and the proposed EMS is applied for day ahead scheduling considering the 3rd scenario. AFS algorithm which is used as a tool to solve the optimization part, has estimated the total generation cost over a day as 3228.9 Euros as against AIMD algorithm, which estimated the same as 3762

Euros. The AFS algorithm is better optimizing the generation mix with a net saving of 16.5%. The proposed EMS able to balance the generation with load at each interval. The main objective of the proposed EMS is reducing the effect of uncertainty of renewable energy generation and it is well attained using battery storage. The EMS is able to schedule battery storage during excess generation and discharge the battery during deficit generation while maintaining the rate of power exchange and state of charge limits. The EMS is also able to implement load shedding when there is no possibility to supply the demand in full during the peak hours, i.e., between 14th to 21st hours.

CONFLICTS OF INTEREST

No conflict of interests was declared by the authors.

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APPENDIX

Abbreviations	
$P_i^{con}(t)$	Power generated by i^{th} generator at time interval t .
$P_j^{res}(t)$	Power generated by j^{th} RES generator at time interval t .
$P_i^{con_{max}}$, $P_i^{con_{min}}$	Maximum and Minimum power limits of i^{th} generator at time interval t .
$P_j^{res_{max}}$, $P_j^{res_{min}}$	Maximum and Minimum power limits of j^{th} RES generator at time interval t .
$W_i^{con}(t), W_j^{res}(t)$	State vectors denoting ON/OFF states of i^{th} generator and j^{th} RES generator at time interval <i>t</i> respectively.
$P_k^{stg_{ch_{max}}}, P_k^{stg_{dch_{max}}}$	Maximum Charge/Discharge rates of k^{th} storage device.
$SoC_k(t)$	State of charge (SoC) of k^{th} storage device at time interval t .
SoC_k^{max} , SoC_k^{min}	Maximum and minimum limits of SoC of k^{th} storage device.
η_k	Efficiency (Discharging/Charging) of k^{th} storage device.
$L_m^{shd}(t)$	m^{th} load shed at time interval t.
$C_m^{shd}(t)$	Compensation due to shedding of m^{th} load at time interval t .
$Y_m^{shd}(t)$	State of vector denoting Connected/Shed state of m^{th} load at time interval t.
$\mathcal{C}_i^{con}(t)$	Power generation cost of i^{th} generator at time interval t .
$C_j^{res}(t)$	Power generation cost of j^{th} RES generator at time interval t .
$c_i^{con}(t)$	Cost/unit of power generation of i^{th} conventional generator at time t.
$C_i^{con_{start}}$, $C_i^{con_{shut}}$	Start-up and shut-down cost of i^{th} conventional generator.
$U_i^{con}(t), V_i^{con}(t)$	State vectors denoting the shut-down/start-up states of i^{th} generator at time interval t.
$U_j^{res}(t), V_j^{res}(t)$	State vectors denoting the shut-down/start-up states of j^{th} renewable generator at time interval t .
$X_i^{con}(t), X_j^{res}(t)$	State vectors denoting the OFF/ON position of the generators and RES at time t .
$C_k^{stg}(t)$	Expenses of energy supply by k^{th} storage device at time interval t .
$c_k^{stg}(t)$	Unit cost energy supply from k^{th} storage device at time interval t .
$Y_k^{stg}(t)$	State vector denoting the type of power exchange by k^{th} storage device at time interval t
$C^{grd}(t)$	Power drawing expenses from grid at time interval t .
$c^{grd}(t)$	Unit cost of grid power at time interval <i>t</i> .
$P^{grd}(t)$	Power drawn from grid at time interval <i>t</i> .
Ψ_j	Vector denoting all numeric values of $P_j^{res}(t)$.
Ψ	Uncertainty set.

$P_{j,s}^{restotal_{min}}, P_{j,s}^{restotal_{max}}$	Upper and lower limits for total energy availability with the j^{th} RES over a time sub- horizon S.
t	Time interval.
$P_i(t)$	Power generated by i^{th} generator at time interval t .
d(t)	Demand at time interval <i>t</i> .
$P_{i_{max}}$, $P_{i_{min}}$	Generation limits of i^{th} generator.