# Skew $\lambda$-cyclic Codes over $Y_{2}$ 

Abdullah DERTL亡 ${ }^{1 *}{ }^{*}$, Yasemin ÇENGELLENMIST룰<br>${ }^{1}$ Ondokuz Mayıs University, Faculty of Arts and Sciences, Mathematics Department, Samsun, Turkey<br>${ }^{2}$ Trakya University, Faculty of Sciences, Mathematics Department, Edirne, Turkey

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#### Abstract

In the present paper, by defining two non-trivial automorphisms and Gray maps over $Y_{2}=\mathbb{Z}_{4}+u \mathbb{Z}_{4}+v \mathbb{Z}_{4}+u v \mathbb{Z}_{4}$, where $u^{2}=0, v^{2}=0, u v=v u$, the algebraic structcure of the skew $\lambda$-cyclic codes and their Gray images over the finite ring $Y_{2}$ are determined, where $\lambda=1+u+v+u v$.


Keywords: Skew codes, Gray map, Finite rings.

## $Y_{2}$ Halkası Üzerinde Skew $\lambda$-cyclic Kodlar

## Öz

Bu çalışmada, $u^{2}=0, v^{2}=0, u v=v u, \lambda=1+u+v+u v$, olmak üzere $Y_{2}=\mathbb{Z}_{4}+u \mathbb{Z}_{4}+v \mathbb{Z}_{4}+u v \mathbb{Z}_{4}$, sonlu halkası üzerinde iki farklı Gray dönüşümü ve otomorfizma tanımlanarak skew $\lambda$-cyclic kodların cebirsel yapısı ve bu kodların Gray görüntüleri belirlenmiştir.

Anahtar Kelimeler: Skew kodlar, Gray dönüşümü, Sonlu halkalar.

## 1. Introduction

Cyclic codes are the most studied class of linear codes with algebraic structure.

Skew polynomial rings form an important family of non-commutative rings. There are many applications in the construction of algebraic codes. As polynomials in skew polynomial ring exhibit many factorizations, there are many more ideals in a skew polynomial ring than in the commutative case. So the researchers on codes have result in the discovery of many new codes with

[^0]better Hamming distance.
Recently, Delphine Boucher et al. gave skew cyclic and skew $\lambda$-cyclic codes defined by using the skew polynomial rings with a nontrivial automorphism, which are generalization of the notion cyclic and constacyclic codes, respectively (Boucher et al., 2007; Boucher et al., 2008).
T. Abualrub, P. Seneviratre studied skew cyclic codes over $F_{2}+v F_{2}$, where $v^{2}=v$ (Abualrub and Seneviratne, 2012). T. Abualrub, A. Ghrayeb, N. Aydın, I. Siap
introduced skew quasi-cyclic codes. They obtained several new codes with Hamming distance exceeding the distance of the previously best known linear codes with comparable parameters (Abualrub et al., 2010). In (Siap et al., 2011), they studied a special type of linear codes called skew cyclic codes in the most general case. M. Bhaintwal studied skew quasi-cyclic codes over Galois rings (Bhaintwal, 2012). Wu investigated the structures of skew cyclic and skew quasi-cyclic of arbitrary length over Galois rings. They shown that the skew cyclic codes are equivalent to either cyclic and quasi-cyclic codes over Galois rings. Moreover, they gave a necessary and sufficient condition for skew cyclic codes over Galois rings to be free ( $\mathrm{Wu}, 2013$ ).

Dertli et al. studied skew codes over the finite ring to increase the probability of obtaining the large minimum distance (Dertli et al., 2015).

In the present paper, the skew $\lambda$-cyclic codes over the finite ring $Y_{2}$ are studied by using two different non-trivial automorphism over $Y_{2}$, which is motivated by the previous works.

## 2. Material and Methods

The ring

$$
\begin{aligned}
& Y_{2}=\mathbb{Z}_{4}+u \mathbb{Z}_{4}+v \mathbb{Z}_{4}+u v \mathbb{Z}_{4}=\{a+b u+c v+d u v: \\
& \left.a, b, c, d \in \mathbb{Z}_{4}, u^{2}=0, v^{2}=0, u v=v u\right\}
\end{aligned}
$$

is commutative ring with $4^{4}$ elements and characteristic 4. A linear code $\wp$ over $Y_{2}$ of length $n$ is a $Y_{2}$ submodule of $Y_{2}^{n}$. An element of $\wp$ is called a codeword.

We defined two Gray maps as follows

$$
\begin{gathered}
\Psi_{1}: Y_{2} \rightarrow \mathbb{Z}_{4}^{4} \\
a+b u+c v+d u v \mapsto(a, b-a, c-a, d-a)
\end{gathered}
$$

and

$$
\begin{gathered}
\Psi_{2}: Y_{2} \rightarrow \mathbb{Z}_{4}^{8} \\
a+b u+c v+d u v \mapsto(a, 3 a, b, 3 b, c, 3 c, d, 3 d)
\end{gathered}
$$

The Gray map $\Psi_{t}$ can be extended to $Y_{2}^{n}$, naturally, for $t=1,2$.

The Lee weight on $\mathbb{Z}_{4}$, denoted $w_{L}$, is defined as

$$
w_{L}(\eta)=\left\{\begin{array}{cc}
0, & \eta=0 \\
1, & \eta=1 \text { or } 3 \\
2, & \eta=2
\end{array}\right.
$$

For any $\xi \in Y_{2}$, the Lee weight of $\xi$ is defined as

$$
w_{L}(\xi)=w_{L}\left(\Psi_{t}(\xi)\right)=\sum_{i=1}^{r} w_{L}\left(\eta_{i}\right)
$$

where $\Psi_{t}(\xi)=\left(\eta_{1}, \eta_{2}, \ldots, \eta_{r}\right)$,
$\eta_{i} \in \mathbb{Z}_{4}, i=1,2, \ldots, r, t=1,2$. The Lee weight of a vector $e=\left(e_{1}, \ldots, e_{n}\right) \in Y_{2}^{n}$ is defined to be a sum of the Lee weights of its components, that is $w_{L}(e)=\sum_{i=1}^{n} w_{L}\left(e_{i}\right)$. Moreover, for any $e_{1}, e_{2} \in Y_{2}^{n}$, the Lee distance between $e_{1}$ and $e_{2}$ is defined as $d_{L}\left(e_{1}, e_{2}\right)=w_{L}\left(e_{1}-e_{2}\right)$.

Theorem 2.1: The Gray map $\Psi_{t}$ is a linear and distance preserving map, for $t=1,2$.

A code $\wp$ over $Y_{2}$ is a $\lambda$-cyclic code with the property that if $e=\left(e_{0}, e_{1}, \ldots, e_{n-1}\right) \in \wp$ then $v(\wp)=\left(\lambda e_{n-1}, e_{0}, \ldots, e_{n-2}\right) \in \wp$ where $\lambda$ is a unit element of $Y_{2}$. A subset $\wp$ of $Y_{2}^{n}$ is a $\lambda-$ cyclic code of length $n$ if and only if it is
polynomial representation is an ideal of $Y_{2}[x] /\left\langle x^{n}-\lambda\right\rangle$. If $\lambda$ is equal to $1(-1)$, the $\wp$ is called cyclic code (negacyclic code), respectively.

## 3. Research Findings

### 3.1. Skew codes over $Y_{2}$

Let $Y_{2}$ be a finite ring and $\Omega_{i}$ be a non-trivial automorphism over $Y_{2}$ and $\lambda=1+u+v+u v$, $i=1,2$.

Definition 3.1.1: A subset $\wp$ of $Y_{2}^{n}$ is called a skew $\lambda$-cyclic code of length $n$ if $\wp$ satisfies the following conditions, for $i=1,2$,

1) $\wp$ is a submodule of $Y_{2}^{n}$
2) If $c=\left(c_{0}, c_{1}, \ldots, c_{n-1}\right) \in \wp$, then
$\sigma_{\Omega_{i}}(c)=\left(\lambda \Omega_{i}\left(c_{n-1}\right), \Omega_{i}\left(c_{0}\right), \ldots, \Omega_{i}\left(c_{n-2}\right)\right) \in \wp$, where $\sigma_{\Omega_{i}}$ is the skew $\lambda$-cyclic shift operator.

By defining two non-trivial automorphisms over $Y_{2}$ as follows, we can define the skew $\lambda$-cyclic codes over $Y_{2}$.

$$
\begin{gathered}
\Omega_{1}: Y_{2} \rightarrow Y_{2} \\
a+b u+c v+d u v \mapsto a+c u+b v+d u v
\end{gathered}
$$

and

$$
\begin{gathered}
\Omega_{2}: Y_{2} \rightarrow Y_{2} \\
a+b u+c v+d u v \mapsto a-b u-c v+d u v
\end{gathered}
$$

The order of $\Omega_{i}$ is 2 , where $i=1,2$.

The rings

$$
\begin{aligned}
& Y_{2}\left[x, \Omega_{i}\right]=\left\{b_{0}+b_{1} x+\ldots+b_{n-1} x^{n-1}:\right. \\
& \left.b_{j} \in Y_{2}, n \in N, j=0,1, \ldots, n-1\right\}
\end{aligned}
$$

are called skew polynomial rings with the usual addition of polynomials and the multiplication as follows

$$
\left(a x^{s}\right)\left(b x^{l}\right)=a \Omega_{i}^{s}(b) x^{s+l}
$$

where $i=1,2$. They are non-commutative rings.

In polynomial representation, a skew $\lambda$ cyclic code of length $n$ over $Y_{2}$ is defined as a left ideal of the quotient ring $A_{\Omega_{i}, n}=Y_{2}\left[x, \Omega_{i}\right] /\left\langle x^{n}-\lambda\right\rangle$, if the order of $\Omega_{i}$ divides $n$, that is $n$ is even. If the order of $\Omega_{i}$ does not divides $n$, a skew $\lambda$-cyclic code of length $n$ over $Y_{2}$ is defined as a left $Y_{2}\left[x, \Omega_{i}\right]$-submodule of $A_{\Omega_{i}, n}$, since the set

$$
\begin{aligned}
& A_{\Omega_{i}, n}=Y_{2}\left[x, \Omega_{i}\right] /\left\langle x^{n}-\lambda\right\rangle= \\
& \left\{f_{i}(x)+\left\langle x^{n}-\lambda\right\rangle: f_{i}(x) \in Y_{2}\left[x, \Omega_{i}\right]\right\}
\end{aligned}
$$

is a left $Y_{2}\left[x, \Omega_{i}\right]$-module, for $i=1,2$.
In both case, the following is hold.
Theorem 3.1.2: Let $\wp$ be a skew $\lambda$-cyclic code over $Y_{2}$ and let $g(x)$ be a polynomial in $\wp$ of minimal degree. If the leading coefficient of $g(x)$ is a unit in $Y_{2}$, then $\wp=\langle g(x)\rangle$, where $g(x)$ is a right divisor of $x^{n}-\lambda$.

Proof: It is proved as in the proof of Lemma 3 and Theorem 1 in (Gao et al., 2017).

Proposition 3.1.3: Let $\Omega_{1}$, vand $\Psi_{1}$ be as above. Then $\Psi_{1} \sigma_{\Omega_{1}}=\rho \Psi_{1} v$, where $\rho$ is a permutation defined by

$$
\rho(x, y, k, p)=(x, k, y, p)
$$

for $x, y, k, p \in \mathbb{Z}_{4}^{n}$.
Proof: Let $e_{i}=a_{i}+b_{i} u+c_{i} v+d_{i} u v$ be the elements of $Y_{2}$ for $i=0,1, \ldots, n-1$. Then

$$
\begin{aligned}
\sigma_{\Omega_{1}}\left(e_{0}, \ldots, e_{n-1}\right) & =\left(\lambda \Omega_{1}\left(e_{n-1}\right), \Omega_{1}\left(e_{0}\right), \ldots, \Omega_{1}\left(e_{n-2}\right)\right) \\
& =\left(\begin{array}{c}
a_{n-1}+\left(a_{n-1}+c_{n-1}\right) u+\left(a_{n-1}+b_{n-1}\right) v+ \\
\left(a_{n-1}+b_{n-1}+c_{n-1}+d_{n-1}\right) u v, a_{0}+c_{0} u+b_{0} v+d_{0} u v \\
\ldots, \ldots, a_{n-2}+c_{n-2} u+b_{n-2} v+d_{n-2} u v
\end{array}\right) .
\end{aligned}
$$

By applying $\Psi_{1}$, we have

$$
\Psi_{1}\left(\sigma_{\Omega_{1}}\left(e_{0}, \ldots, e_{n-1}\right)\right)=\binom{a_{n-1}, \ldots a_{n-2}, c_{n-1}, \ldots, c_{n-2}-a_{n-2}, b_{n-1}, \ldots,}{b_{n-2},-a_{n-2}, b_{n-1}+c_{n-1}+d_{n-1}, \ldots, d_{n-2}-a_{n-2}} .
$$

On the other hand

$$
\Psi_{1} v\left(e_{0}, \ldots, e_{n-1}\right)=\binom{a_{n-1}, \ldots a_{n-2}, b_{n-1}, \ldots, b_{n-2}-a_{n-2}, c_{n-1}, \ldots,}{c_{n-2},-a_{n-2}, b_{n-1}+c_{n-1}+d_{n-1}, \ldots, d_{n-2}-a_{n-2}} .
$$

If we apply $\rho$, we have

$$
\rho \Psi, v\left(e_{0}, \ldots, e_{n-1}\right)=\binom{a_{n-1}, \ldots a_{n-2}, c_{n-1}, \ldots, c_{n-2}-a_{n-2}, b_{n-1}, \ldots,}{b_{n-2},-a_{n-2}, b_{n-1}+c_{n-1}+d_{n-1}, \ldots, d_{n-2}-a_{n-2}}
$$

We have the expected result.
Theorem 3.1.4: The Gray image of a skew $\lambda$-cyclic code over $Y_{2}$ of length $n$ is permutation equivalent to a $\lambda$-cyclic code over $\mathbb{Z}_{4}$ of length $4 n$.

Proof: Let $\wp$ be a skew $\lambda$-cyclic code over $Y_{2}$ of length $n$. That is $\sigma_{\Omega_{1}}(\wp)=\wp$. If we apply $\quad \Psi_{1}$, we have $\Psi_{1}\left(\sigma_{\Omega_{1}}(\wp)\right)=\Psi_{1}(\wp)$.

From Proposition 3.1.3, we get $\Psi_{1}\left(\sigma_{\Omega_{1}}(\wp)\right)=\Psi_{1}(\wp)=\rho \Psi_{1} v(\wp)$. So $\Psi_{1}(\wp)$ is permutation equivalent to a $\lambda$-cyclic code over $\mathbb{Z}_{4}$ of length $4 n$.

Proposition 3.1.5: Let $\Omega_{2}, v$ and $\Psi_{2}$ be as above. Then $\Psi_{2} \sigma_{\Omega_{2}}=\psi \Psi_{2} v$, where $\psi$ is a permutation defined by

$$
\begin{aligned}
& \psi\left(x_{1}, x_{2}, x_{3}, x_{4}, x_{5}, x_{6}, x_{7}, x_{8}\right)= \\
& \left(x_{1}, x_{2}, x_{4}, x_{3}, x_{6}, x_{5}, x_{7}, x_{8}\right) \\
& \text { for } x_{i} \in \mathbb{Z}_{4}^{n}, i=1,2, \ldots, 8
\end{aligned}
$$

Proof: It is proved as in the proof of the Proposition 3.1.3.

Theorem 3.1.6: The Gray image of a skew $\lambda$-cyclic code over $Y_{2}$ of length $n$ is permutation equivalent to a $\lambda$-cyclic code over $\mathbb{Z}_{4}$ of length $8 n$.

Proof: It is proved as in the proof of the Theorem 3.1.4.

Example 3.1.7: Let $n=3$. We have
$x^{3}-\lambda=\left(x^{2}+(1-u-v+u v) x+(1-u-v)\right)(x-(1-u-v+u v))$ in $Y_{2}\left[x, \Omega_{i}\right]$, for $i=1,2$.

Let $g(x)=x-(1-u-v+u v)$. Then $g(x)$ generates a skew $\lambda$-cyclic code of length 3 with the minimum distance $d=2$. This code is permutation equivalent to $a \lambda$ cyclic code of length $12(24)$ over $\mathbb{Z}_{4}$.

## 4. Conclusion

The skew $\lambda$-cyclic codes over the finite ring $Y_{2}$ are studied because of to increase the probability of obtaining the large minimum
distance. A new two Gray maps and two nontrivial automorphisms over $Y_{2}$ are defined and the Gray images of skew $\lambda$-cyclic codes are determined. So, we can obtain many new codes with better Hamming distance.

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[^0]:    *Corresponding Author:abdullah.dertli@gmail.com

